

# Pointers (1A)

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# Pointers

# Variables

```
int a ;
```

a can hold an integer data

type : int  
size : 4 bytes  
value : integer

```
a = 100 ;
```

a holds the integer 100

address      data

&a            a ← 100

# Pointer Variables

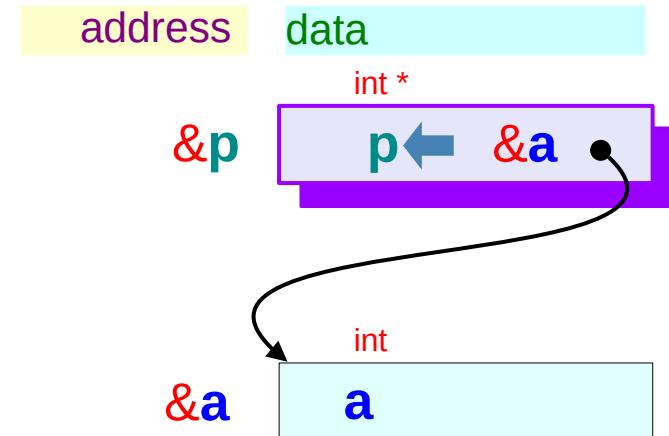
```
int * p ;
```

p can hold an address of  
an integer data

```
p = &a ;
```

p can hold an address (&a)  
of an integer data (a)

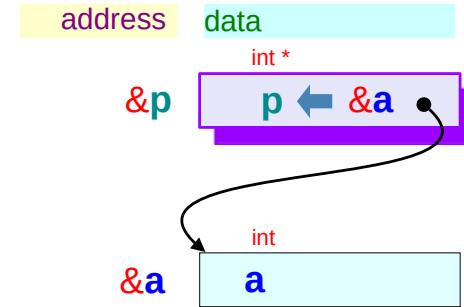
type : int \*  
size : 4 bytes (32-bit system)  
 : 8 bytes (64-bit system)  
value : address



# Dereferenced Pointer variables

```
int * p ;
```

p can hold an address of  
an integer data



*type*            *variable*

int \* **p** ;

*pointer to int*

int \* **p** ;

*int*

address      data

int \*

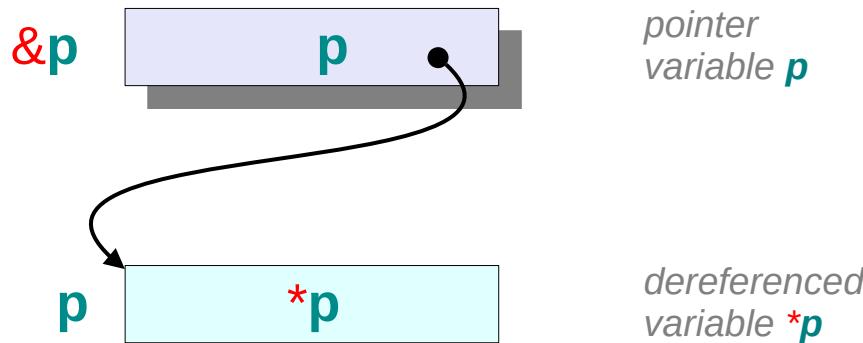
&p      p

int  
\*p ← 100

# Pointer variable **p** and dereferenced variable **\*p**

**&p** : the address of pointer variable **p**

pointer **p** points here,  
thus the address is **p**



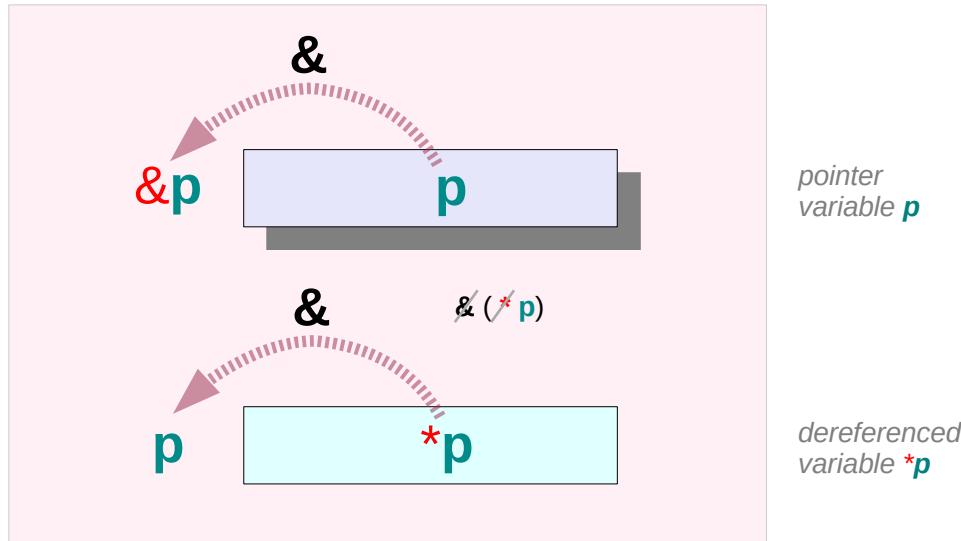
pointer  
variable **p**

derefenced  
variable **\*p**

**\*p** and **p** are variables  
**&p** is an address value

# Address-of operator and dereferencing operator (1)

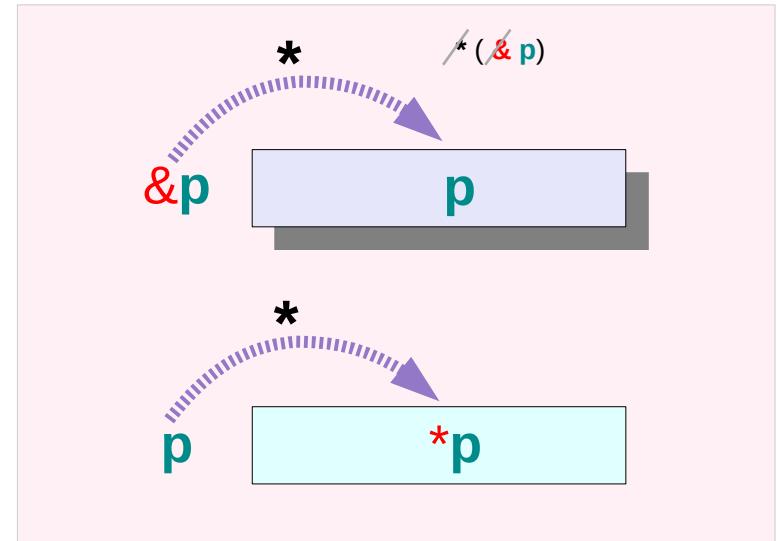
*the address of a variable :  
address-of operator &*



$$\& ( p ) \equiv \text{value}(\&p) \dots \text{value}$$

$$\& (*p) \equiv \text{value}(p) \dots \text{value}$$

*the content at an address :  
derefencing operator \**



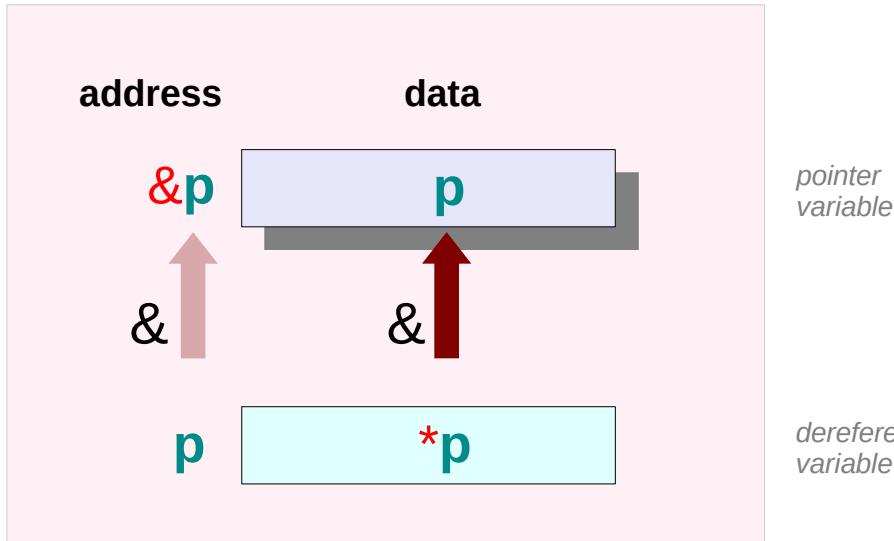
$$* ( \&p ) \equiv p \dots \text{variable}$$

$$* ( p ) \equiv *p \dots \text{variable}$$

# Address-of operator and dereferencing operator (2)

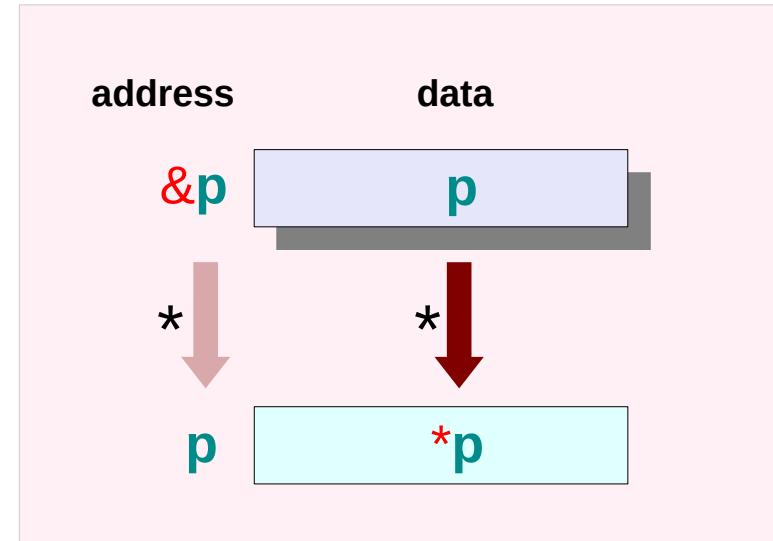
*the address of a variable :  
address-of operator &*

*the content at an address :  
dereferencing operator \**



$\& ( p ) \equiv \text{value}(\&p) \dots \text{value}$

$\& (*p) \equiv \text{value}(p) \dots \text{value}$

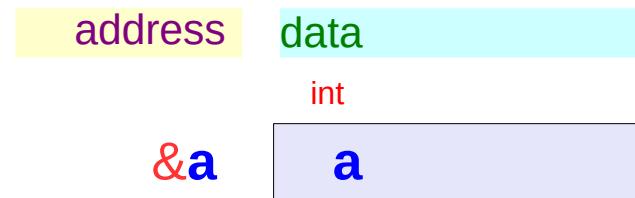


$* (\&p) \equiv p \dots \text{variable}$

$* ( p ) \equiv *p \dots \text{variable}$

# Variables and their addresses

```
int    a ;
```



```
int * p ;
```

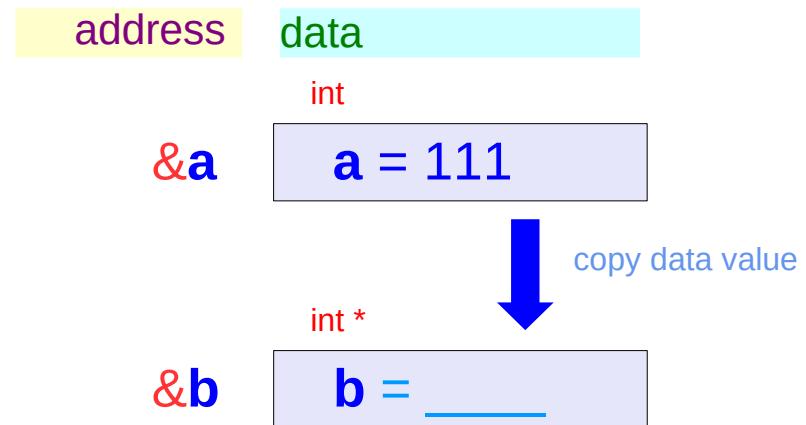


# Assignment of a value

```
int a ;
```

```
int b ;
```

```
b = a ;
```

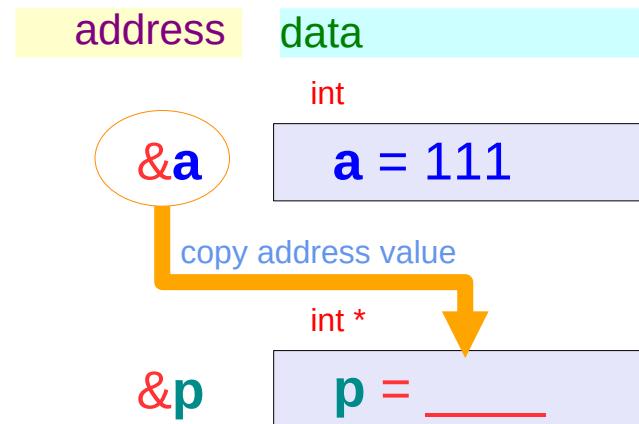


# Assignment of an address value

```
int    a ;
```

```
int * p ;
```

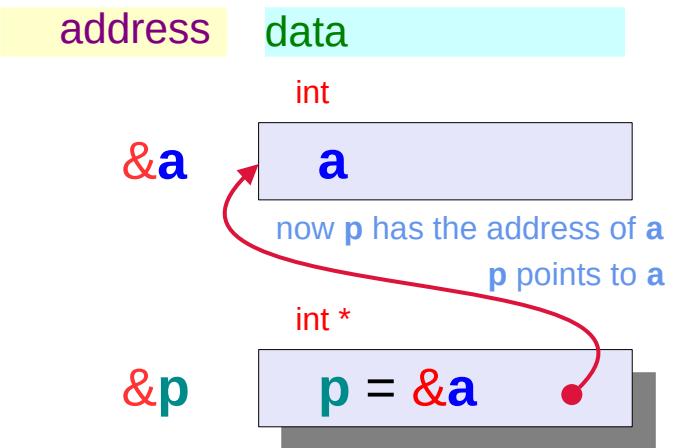
```
p = &a;
```



# Arrow notations

```
int    a ;
```

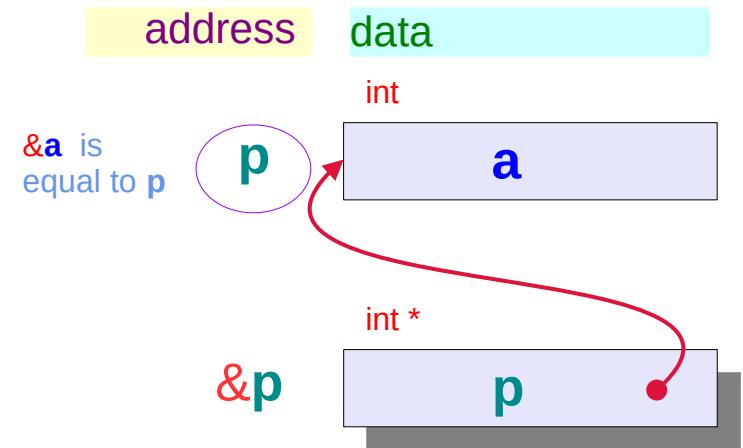
```
int * p = &a;    p = &a;
```



# Pointed address : p

int a;

int \* p = &a;      p = &a;

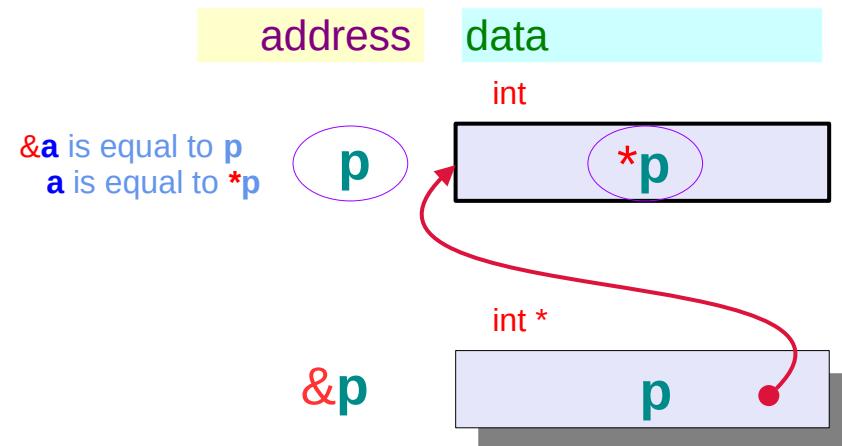


$$p \equiv \&a$$

## Dereferenced variable : \*p

int a;

int \* p = &a;      p = &a;



## equivalence

p     ≡     &a

$$* \text{ (p)} \quad \equiv \quad * \text{ (&a)}$$

$${}^{\textcolor{red}{*}} \textcolor{teal}{p} \equiv \textcolor{blue}{a}$$

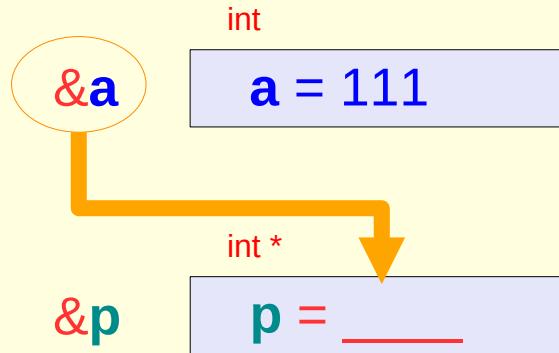
# Variable p, \*p

```
int a;
```

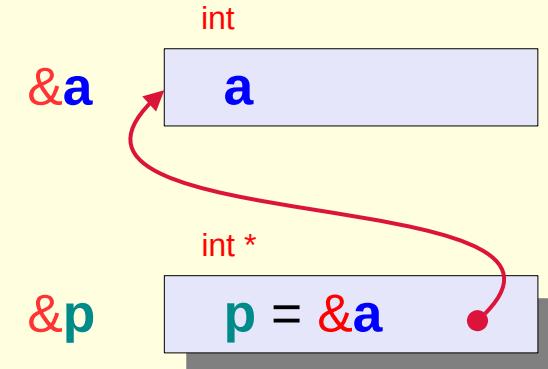
```
int * p;
```

```
p = &a;
```

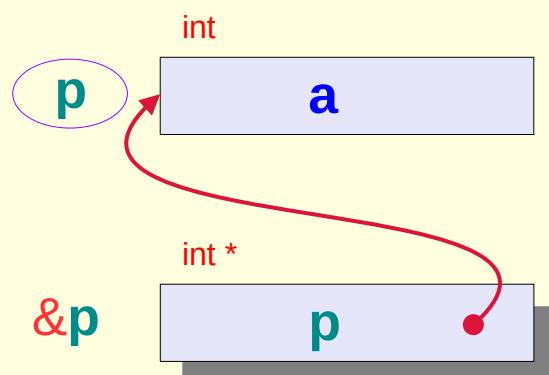
(1) copy address value



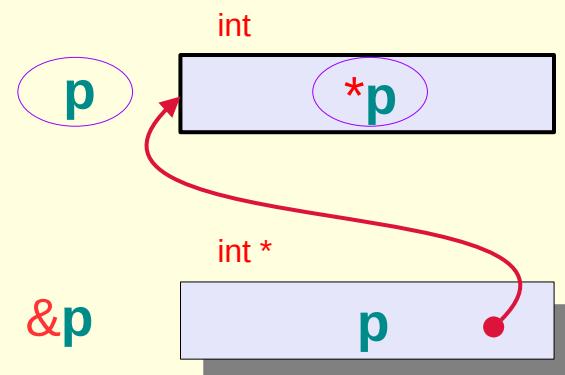
(2) p points to a



(3) &a is equal to p

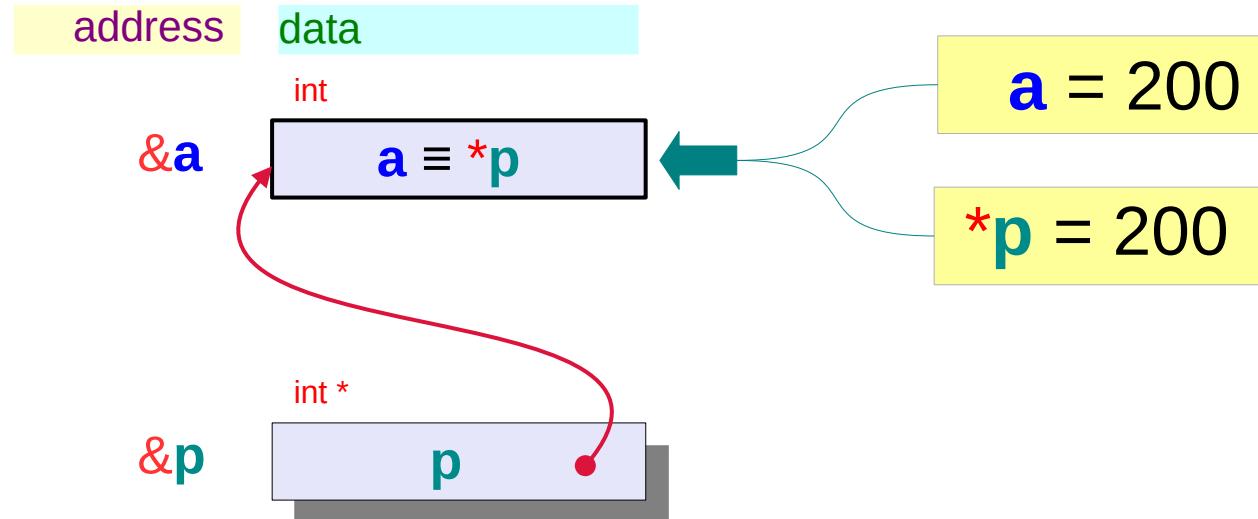


(4) a is equal to \*p



# Two way to access: **a** and **\*p**

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- 
- 1) Read / Write **a**
  - 2) Read / Write **\*p**
-

# Double Pointers

# Double Pointer Variable Definition

```
int ** q ;
```

q holds an address

int \* \*      q ;  
a pointer to  
an integer pointer

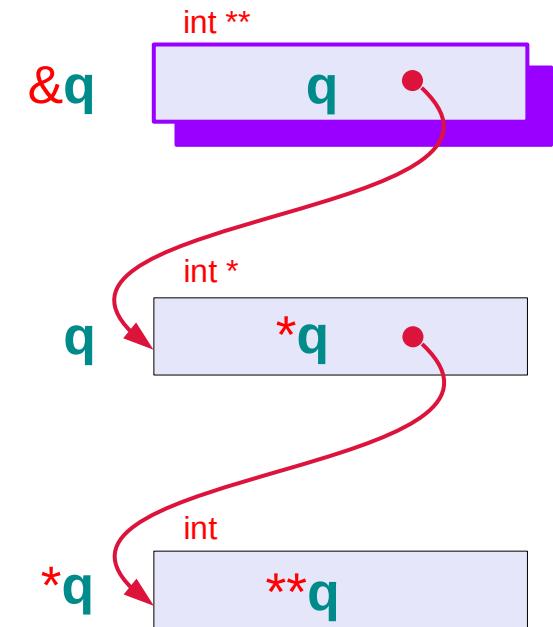
q holds an address of  
an integer pointer data

int \*      \*q ;  
a pointer to  
an integer

\*q holds an address of  
an integer data

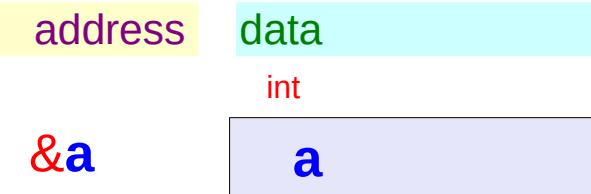
int      \* \*q ;  
an integer

\*\*q holds an integer data

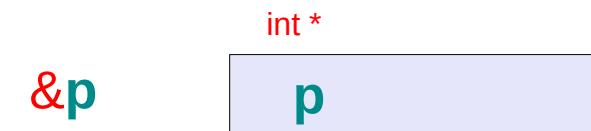


# Variables and their addresses

int a ;



int \* p ;



int \*\* q ;

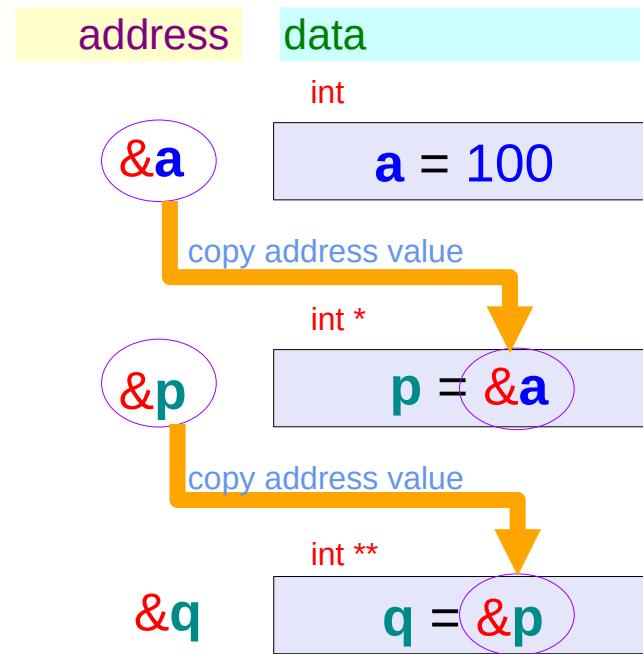


# Assignments of address values

```
int      a = 100 ;
```

```
int * p = &a ;           p = &a;
```

```
int ** q = &p ;         q = &p;
```



# Arrow notations

```
int    a = 100 ;
```

```
int * p = &a ;
```

```
int ** q = &p ;
```

```
p = &a;
```

```
q = &p;
```

address data

int

&a

a = 100

now p has the address of a

p points to a

&p

p = &a

now q has the address of p

q points to p

&q

q = &p

# Pointed addresses : p, q

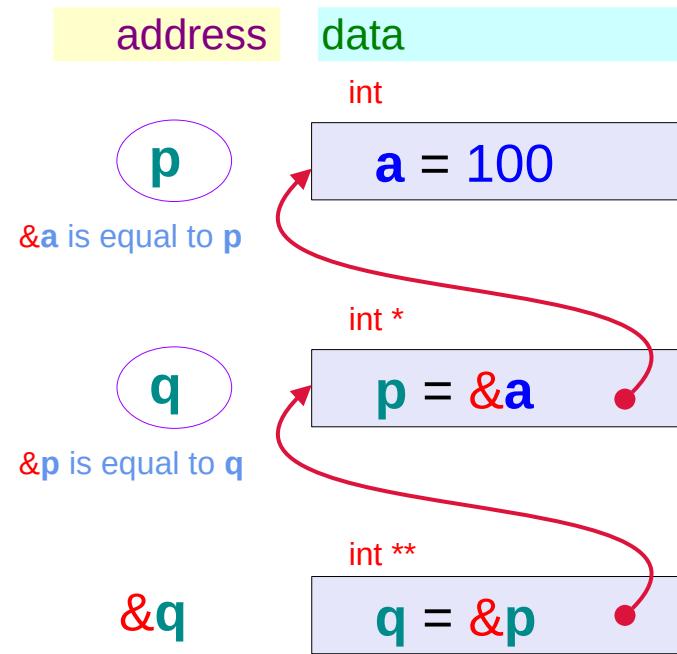
```
int      a = 100 ;
```

```
int * p = &a ;
```

```
int ** q = &p ;
```

**p = &a;**

**q = &p;**



# Dereferenced variables : $\ast q$ , $\ast\ast q$

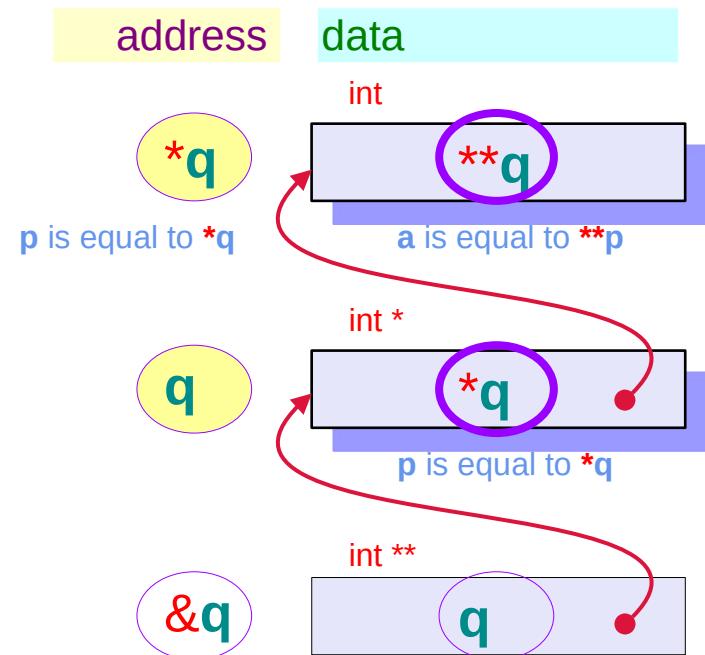
int      a = 100 ;

$$\ast\ast q \equiv a$$

int \* p = &a ;

$$\ast q \equiv p$$

int \*\* q = &p ;

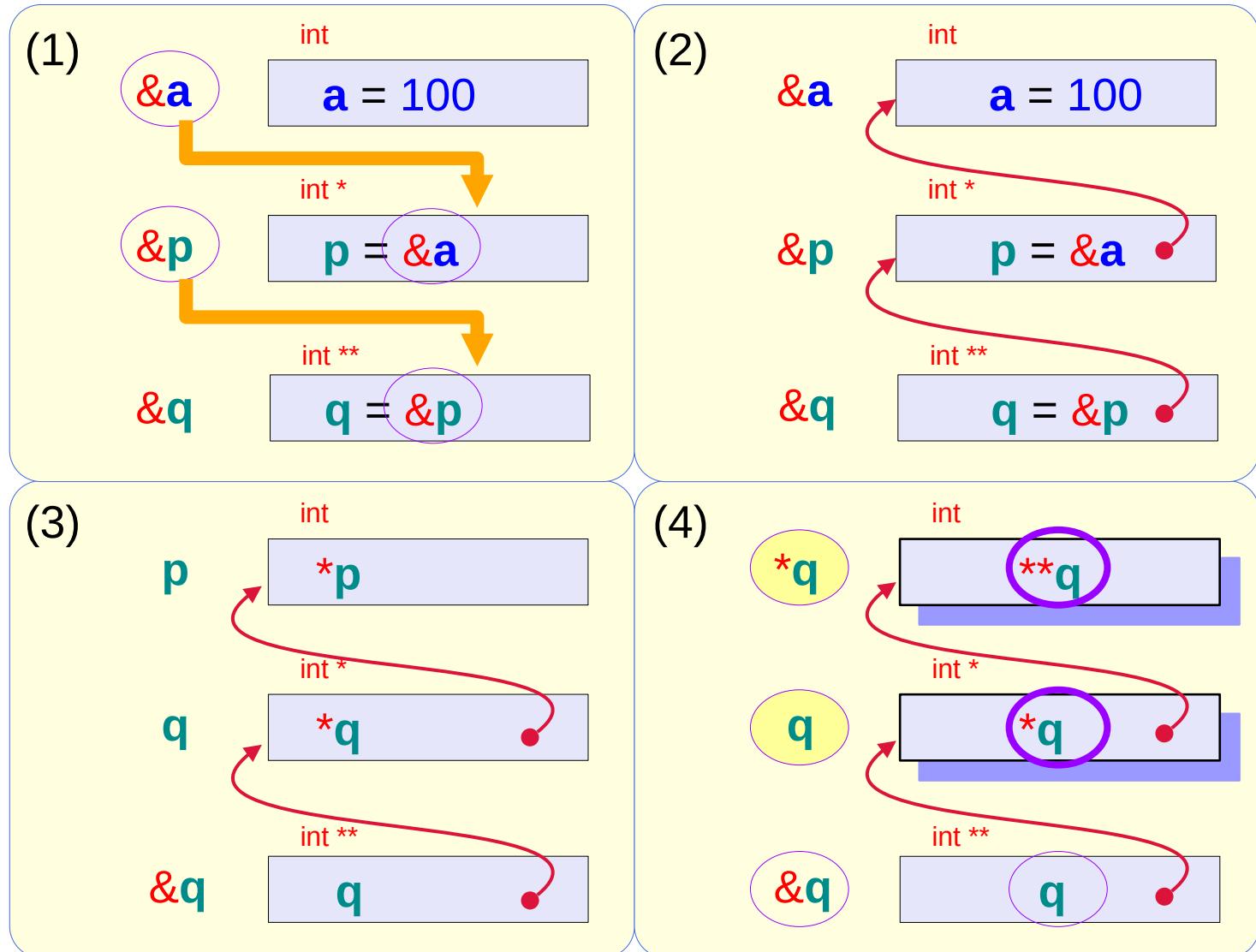


# Variable p, \*p, and q, \*q, \*\*q

```
int a;  
int * p;  
int ** q;
```

**p = &a;**

**q = &p;**



# Aliased variables : \*p,\*q, \*\*q

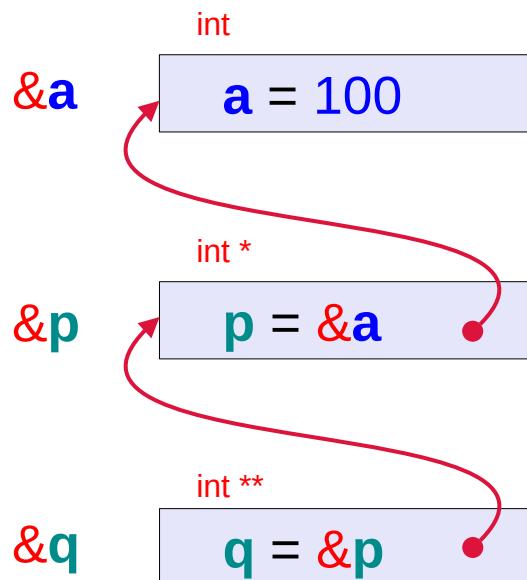
```
int      a = 100 ;  
int  * p = &a ;  
int ** q = &p ;
```

Address  
assignment

$p = \&a ; \rightarrow *p \equiv a$

$p \equiv \&a$   
 $*(p) \equiv *(\&a)$   
 $*p \equiv a$

Variable  
aliasing

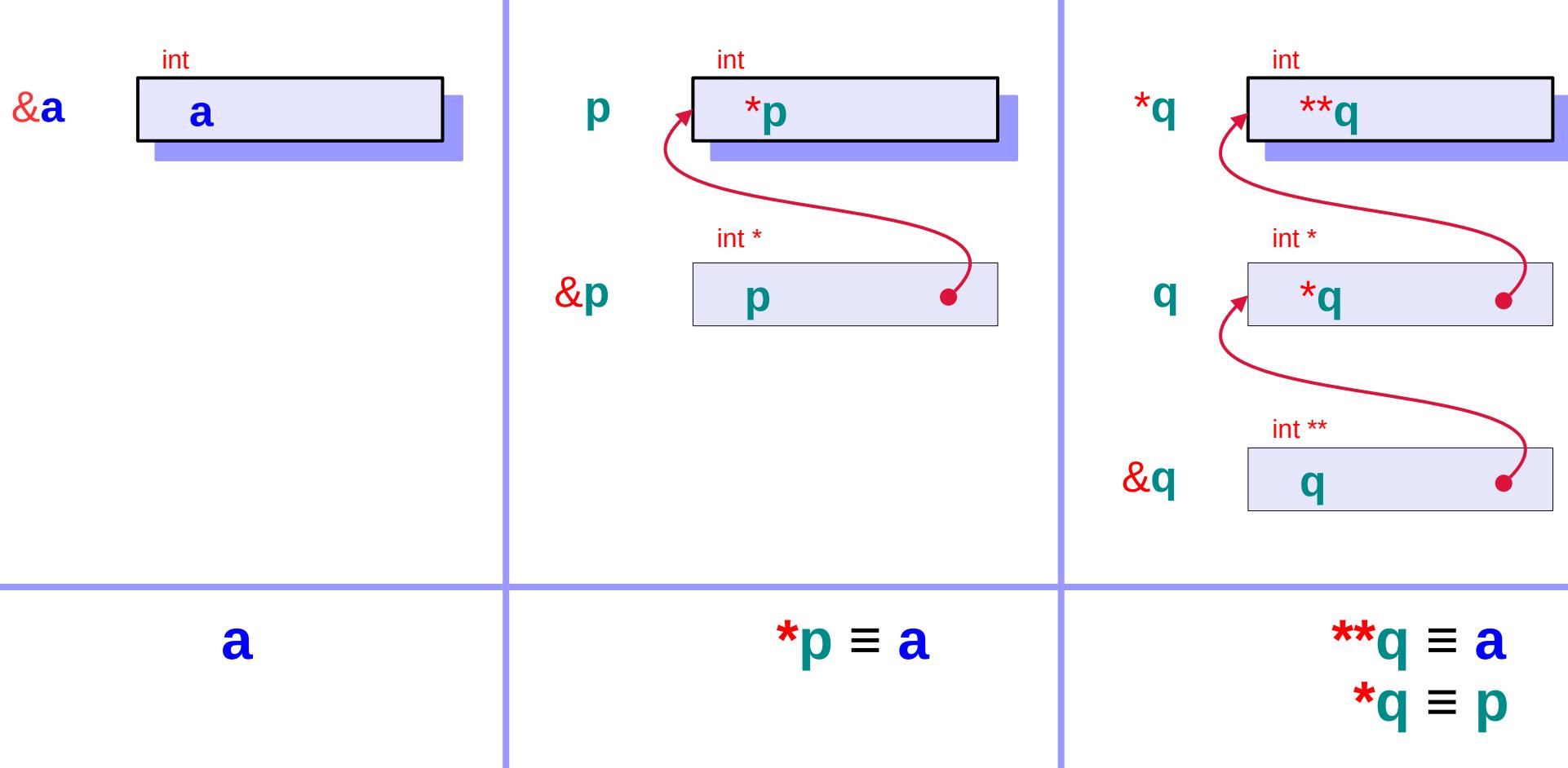


$q = \&p ; \rightarrow *q \equiv p$   
 $\rightarrow **q \equiv a$

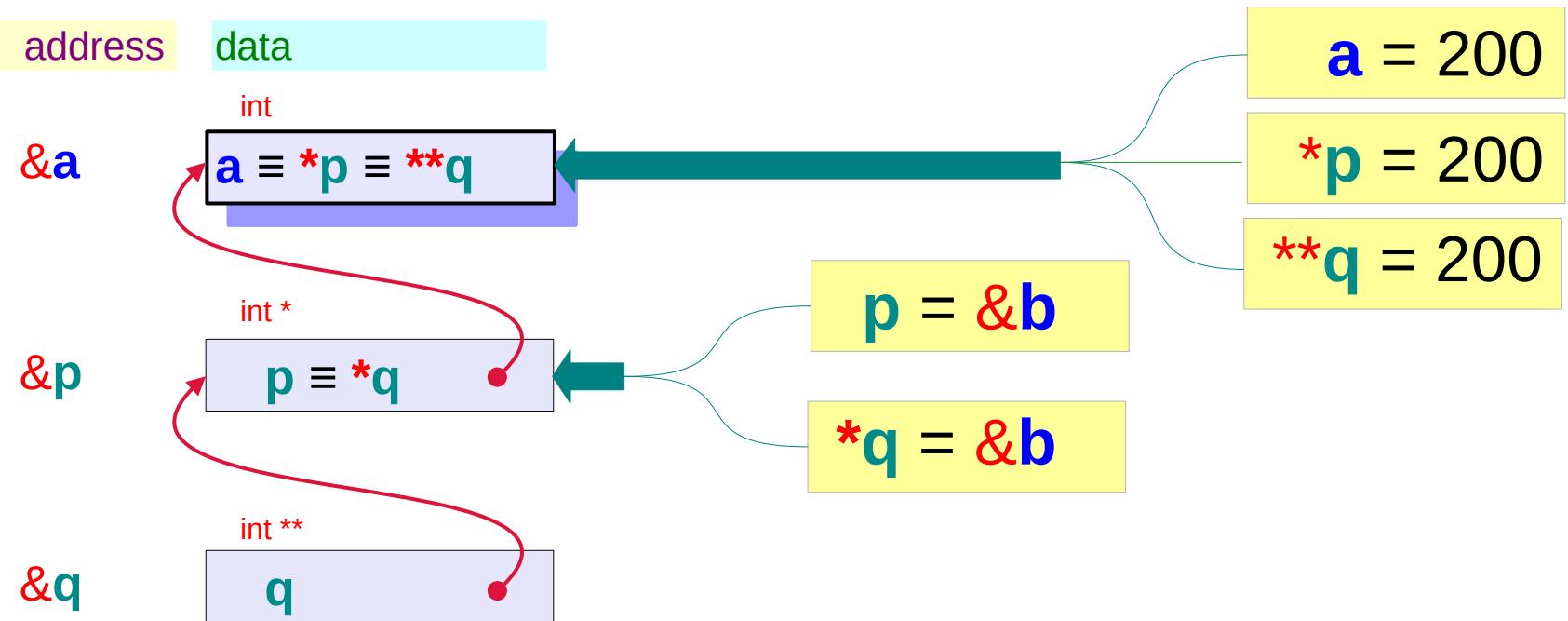
$q \equiv \&p$   
 $*(q) \equiv *(\&p)$   
 $*q \equiv p$   
 $**q \equiv *p$   
 $**q \equiv a$

equivalent relations after  
address assignment

# Two aliased variables of `a : *p, **q`



# Two more ways to access a : \*p, \*\*q



1) Read / Write *p*  
2) Read / Write *\*q*

1) Read / Write *a*  
2) Read / Write *\*p*  
3) Read / Write *\*\*q*

# Single and Double Pointers

# Pointed Addresses and Data

```
int a ;           int  
    &a   a =100
```

The integer variable **a** holds an integer data

```
int * p ;        int *  
    &p   p •          p   int  
                           *p = 200
```

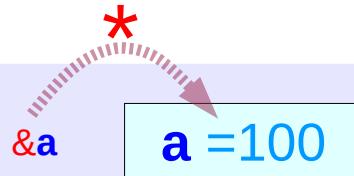
The **pointer** variable **p** holds an **address**,  
at this address, an **integer data** is stored

```
int ** q ;       int **  
    &q   q •          q   int *  
                           *q •          *q   int  
                           **q = 30
```

The **pointer** variable **q** holds an **address**,  
at the address **q**, another **address** **\*q** is stored,  
at the address **\*q**, an **integer data** **\*\*q** is stored

# Dereferencing Operator \*

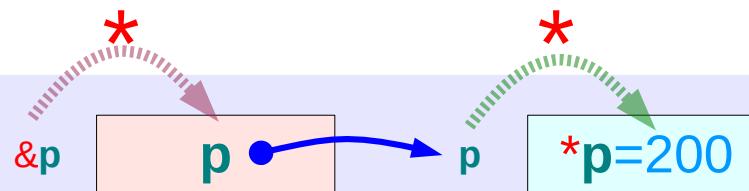
`int a ;`



the expression `a` is a *variable* that can be assigned

$$*(\&a) \equiv a$$

`int * p ;`

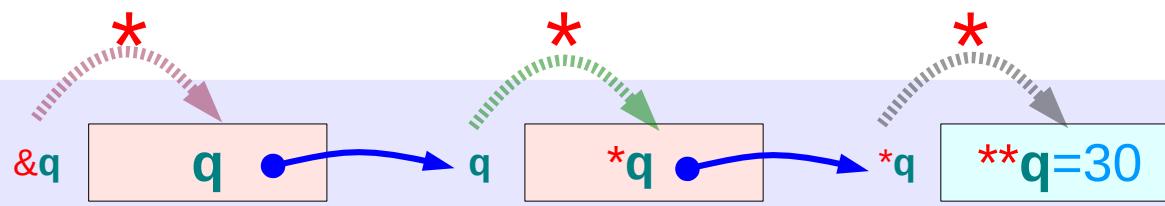


the expression `*p` is a *variable* that can be assigned

$$*(\&p) \equiv p$$

$$*(p) \equiv *p$$

`int ** q ;`



the expression `*q`, `**q` are *variables* that can be assigned

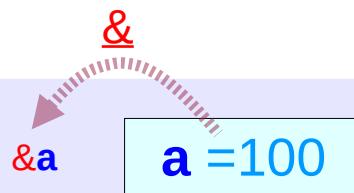
$$*(\&q) \equiv q$$

$$*(q) \equiv *q$$

$$*(\ast q) \equiv \ast\ast q$$

# Math operator & : the inverse operator of \*

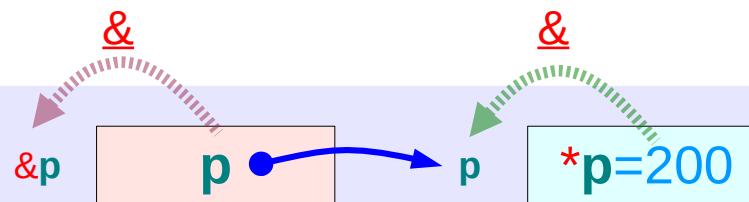
int a ;



\*&x returns x  
&\*y returns y  
& is the inverse operator of \*

int \* p ;

&a

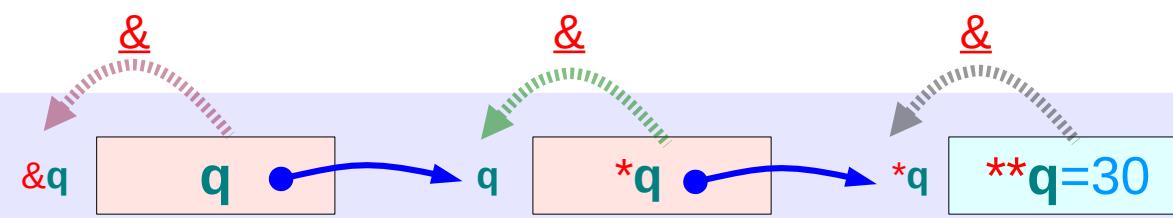


\*&x returns x  
&\*y returns y  
& is the inverse operator of \*

int \* \* q ;

&p

&(\*p) ≡ p



\*&x returns x  
&\*y returns y  
& is the inverse operator of \*

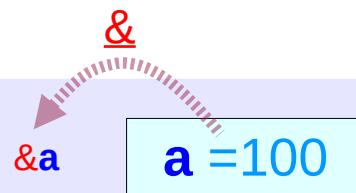
&q

&(\*q) ≡ q

&(\*\*q) ≡ \*q

# C operator & : address-of operator

```
int a ;
```

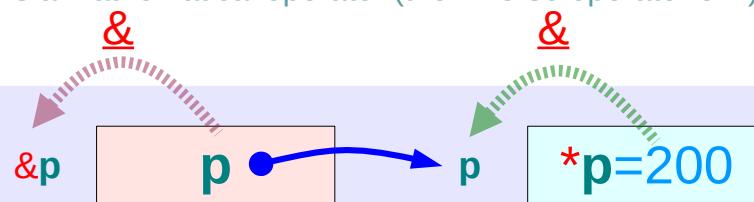


the C expression `&a` returns  
an *address value*  
that cannot be assigned

$$\&a \equiv \text{value}(\underline{\&a})$$

`&` is a mathematical operator (the inverse operator of `*`)

```
int * p ;
```



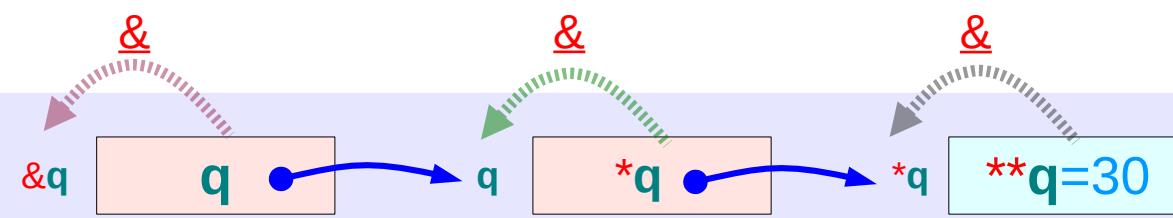
the C expression `&p` returns  
an *address value*  
that cannot be assigned

$$\&p \equiv \text{value}(\underline{\&p})$$

$$\&(*p) \equiv \text{value}(p)$$

`&` is a mathematical operator (the inverse operator of `*`)

```
int ** q ;
```



the C expression `&q` returns  
an *address value*  
that cannot be assigned

$$\&q \equiv \text{value}(\underline{\&q})$$

$$\&(*q) \equiv \text{value}(q)$$

$$\&(**q) \equiv \text{value}(*q)$$

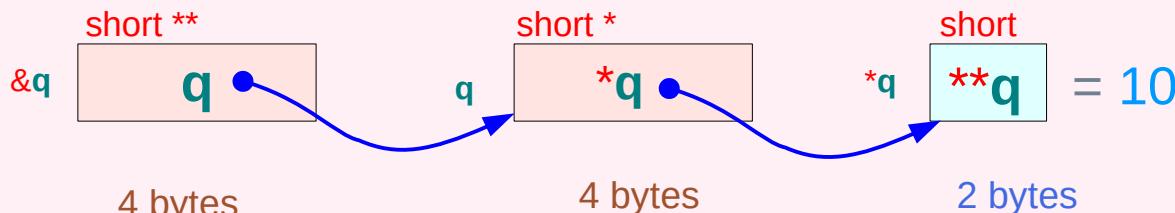
`&` is a mathematical operator (the inverse operator of `*`)

# Pointers to various types on 32-bit system (1)



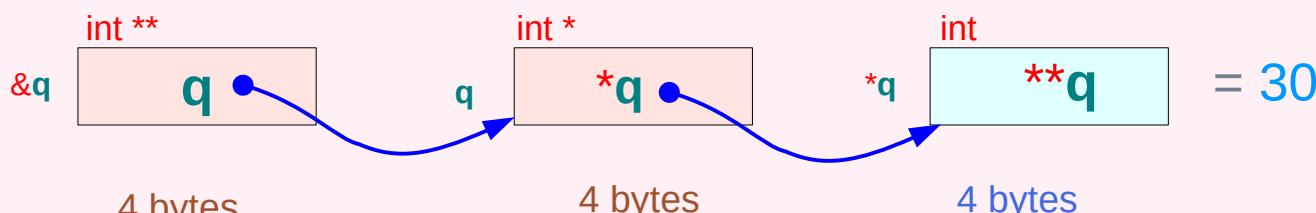
```
char a;  
char * p = &a;  
char ** q = &p;
```

on a 32-bit system



```
short a;  
short * p = &a;  
short ** q = &p;
```

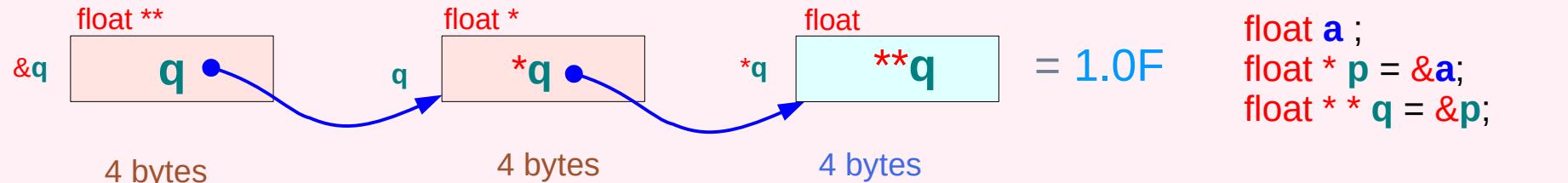
on a 32-bit system



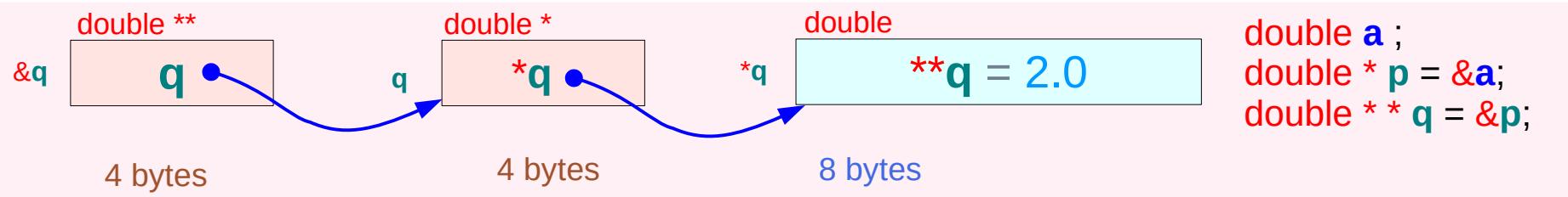
```
int a;  
int * p = &a;  
int ** q = &p;
```

on a 32-bit system

# Pointers to various types on 32-bit system (2)



on a 32-bit system



on a 32-bit system

# Neighborhood of \*p

```
int      a = 100 ;  
int * p = &a ;
```

int \*  
&p      p

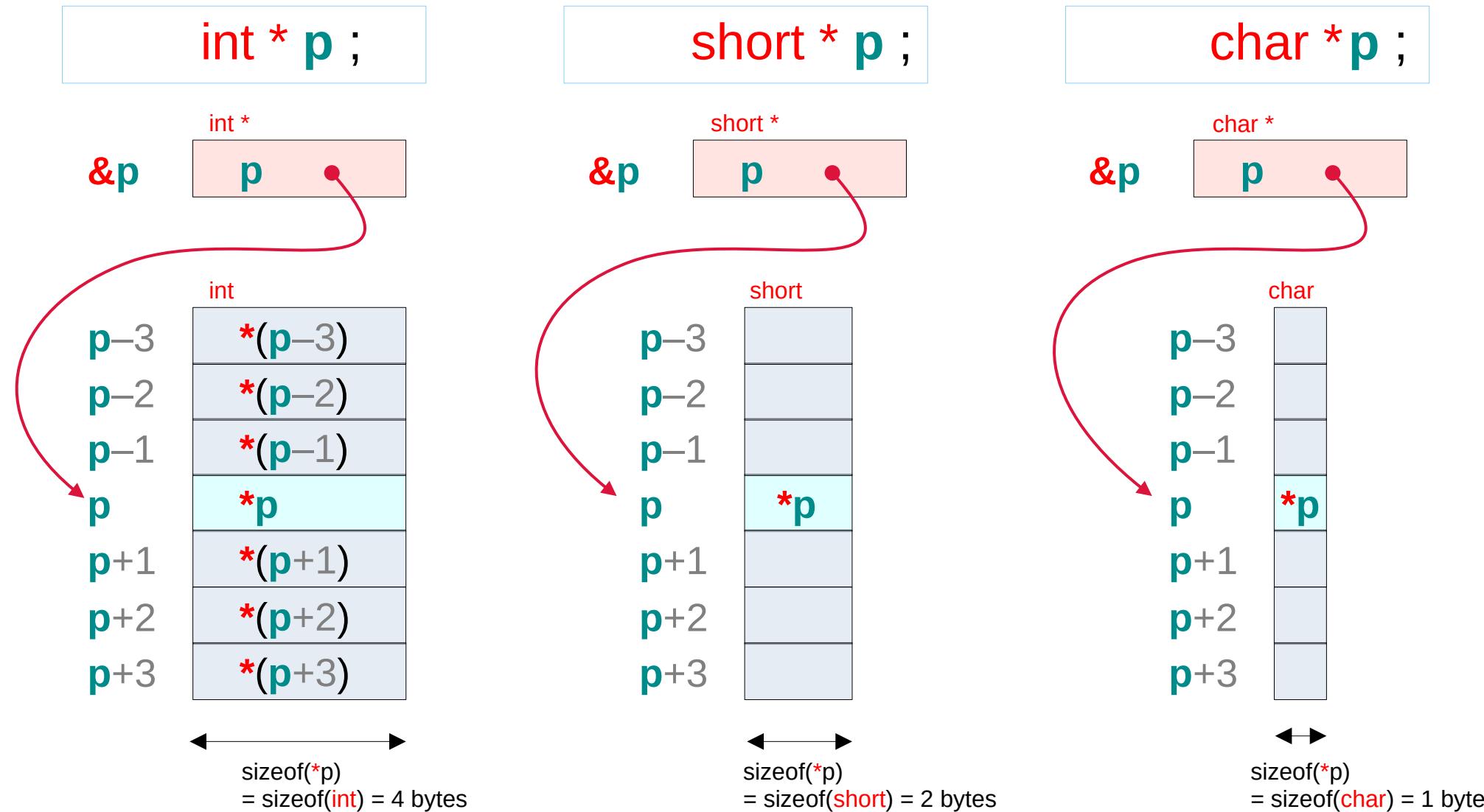
abstract  
address

int \*  
p-3  
p-2  
p-1  
p  
p+1  
p+2  
p+3

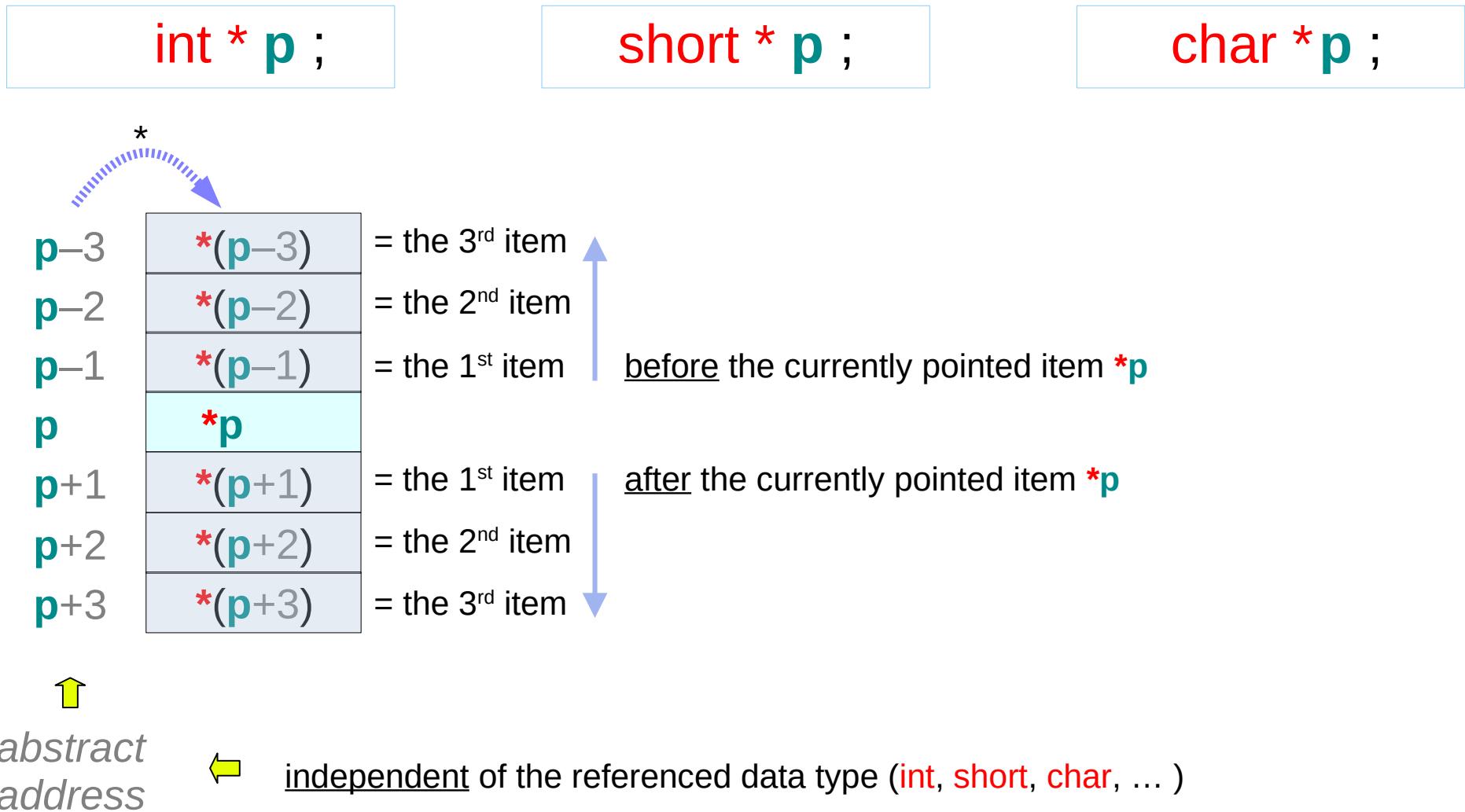
referenced  
data

int	•
int *	•
$\ast(p-3)$	•
$\ast(p-2)$	•
$\ast(p-1)$	•
$\ast p$	•
$\ast(p+1)$	•
$\ast(p+2)$	•
$\ast(p+3)$	•
	•
	•
	•

# Neighborhood of `*p` – various sizes of referenced data



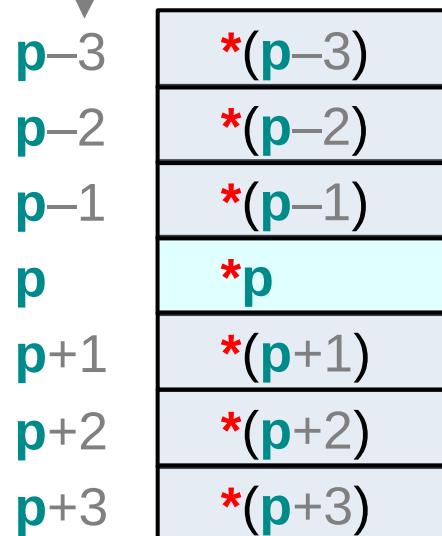
# Neighborhood of `*p` – the same abstract address



# Neighborhood of $*p$ – word aligned view

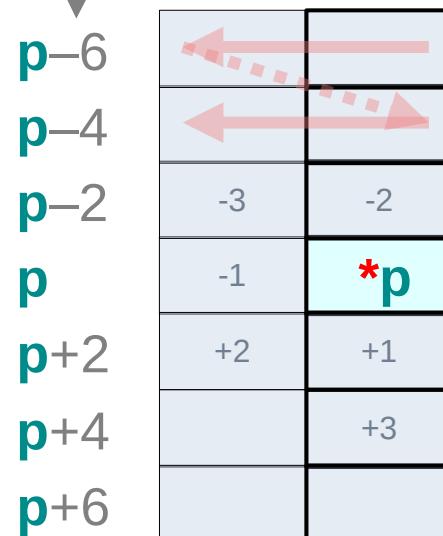
int \* p ;

Word Address



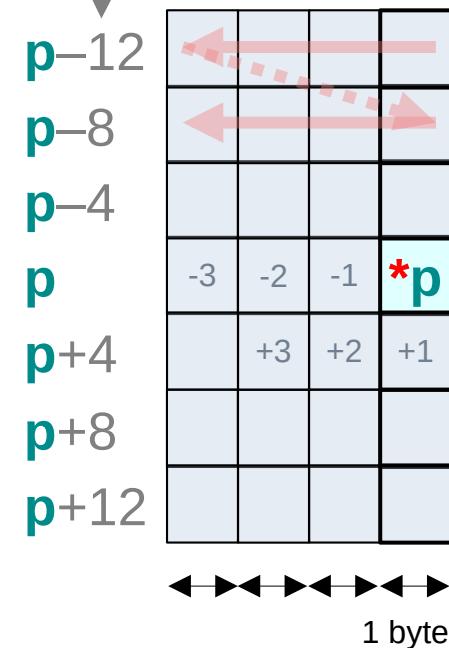
short \* p ;

Least Halfword Address  
among 2 halfwords  
word aligned



char \* p ;

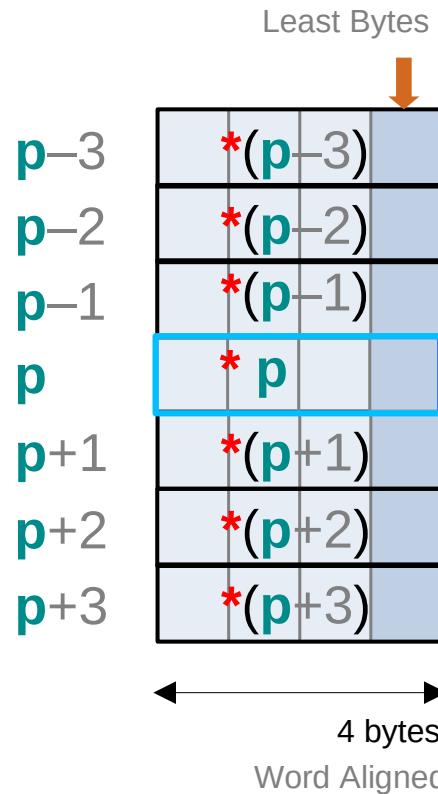
Least Byte Address  
among 4 bytes  
word aligned



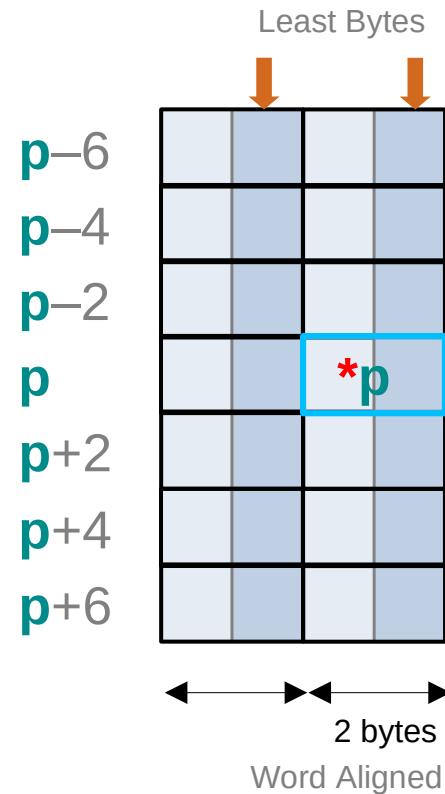
**word view**

# Neighborhood of $*p$ – least bytes

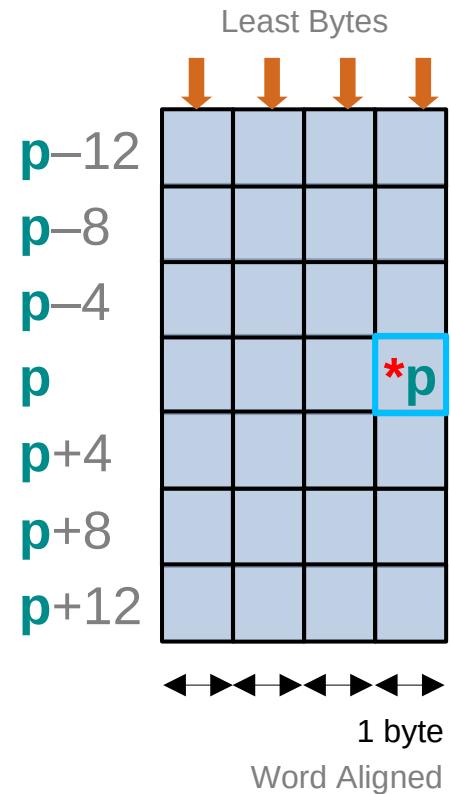
int \* p ;



short \* p ;

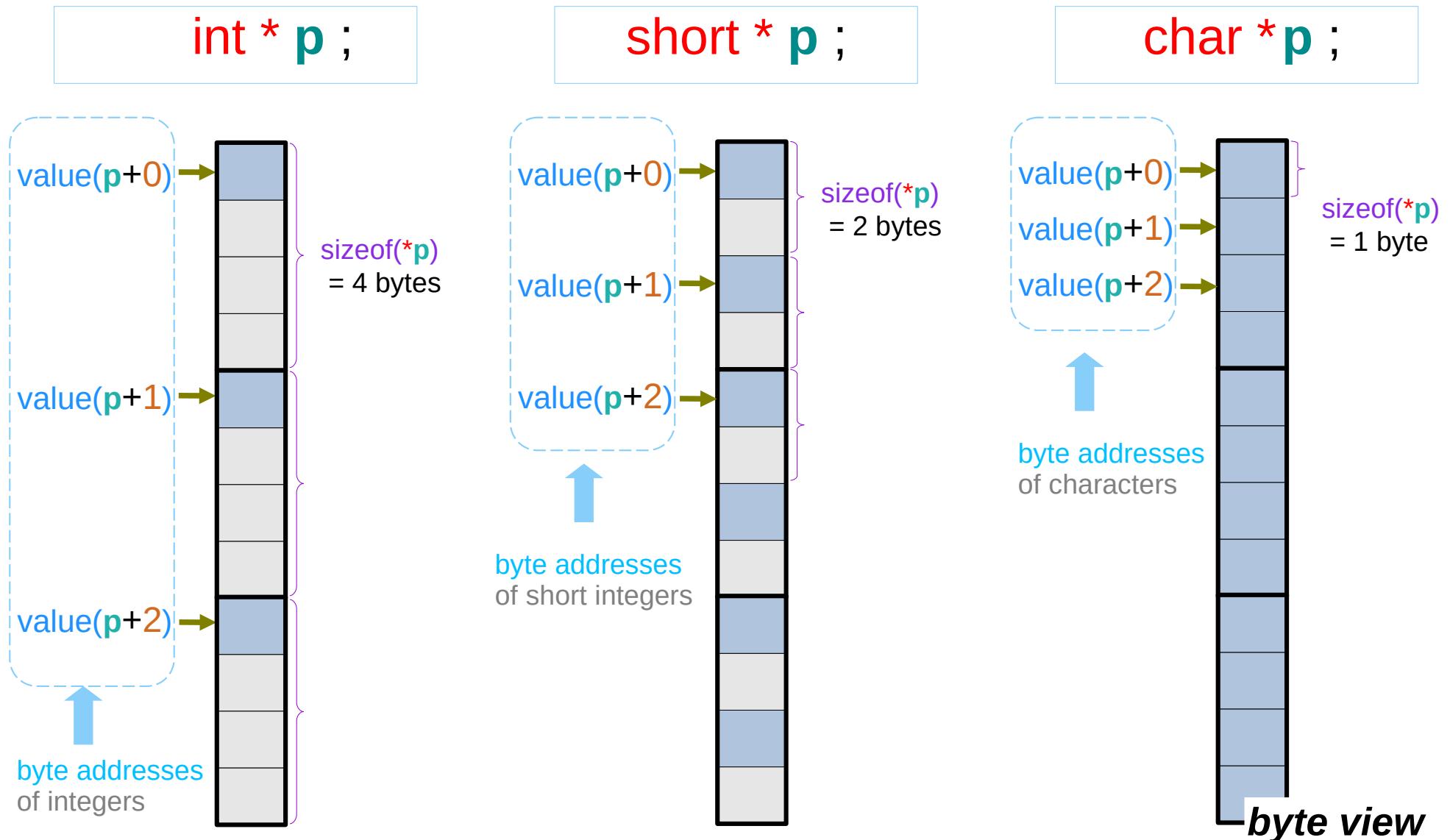


char \* p ;



***word view***

# Neighborhood of $*p$ – least byte addresses



# Neighborhood of $*p$ – incremented pointer values

**int \* p ;**

abstract address  
expression  $p+1$

least byte address  
of int type  $*(p+1)$

$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$

$\text{value}(p+i) = \text{value}(p) + i*4$

**short \* p ;**

abstract address  
expression  $p+1$

least byte address  
of short type  $*(p+1)$

$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$

$\text{value}(p+i) = \text{value}(p) + i*2$

**char \* p ;**

abstract address  
expression  $p+1$

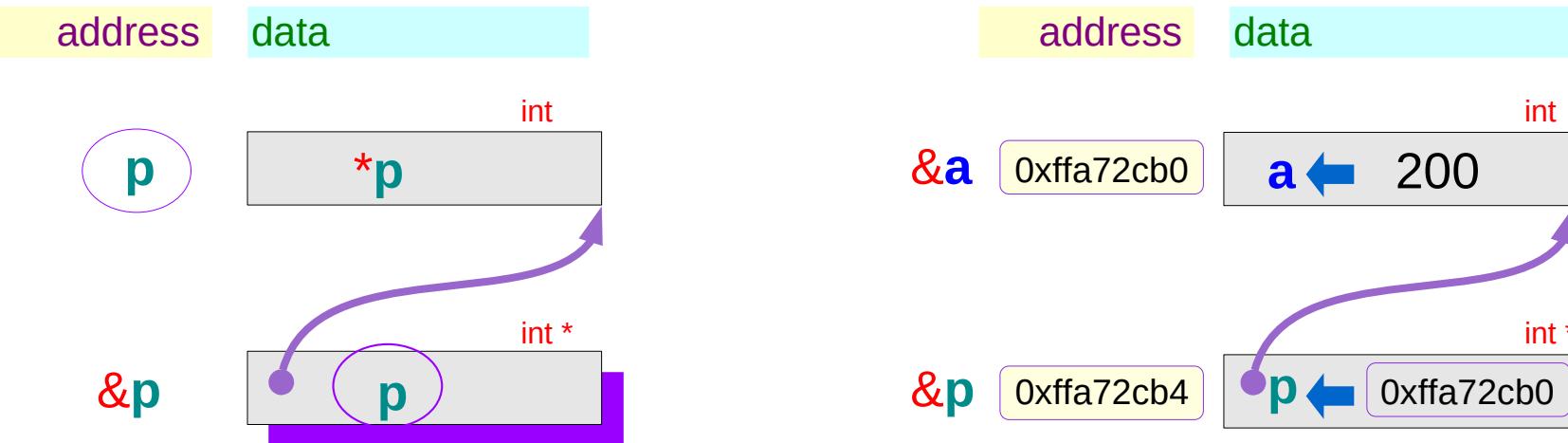
least byte address  
of char type  $*(p+1)$

$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$

$\text{value}(p+i) = \text{value}(p) + i*1$

# Pointer Variable Example

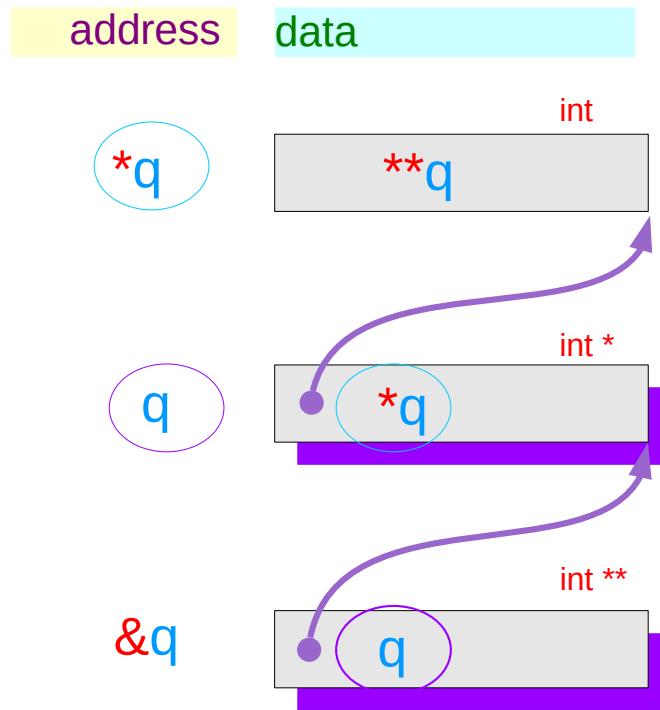
# Pointer variable p of the type int \*



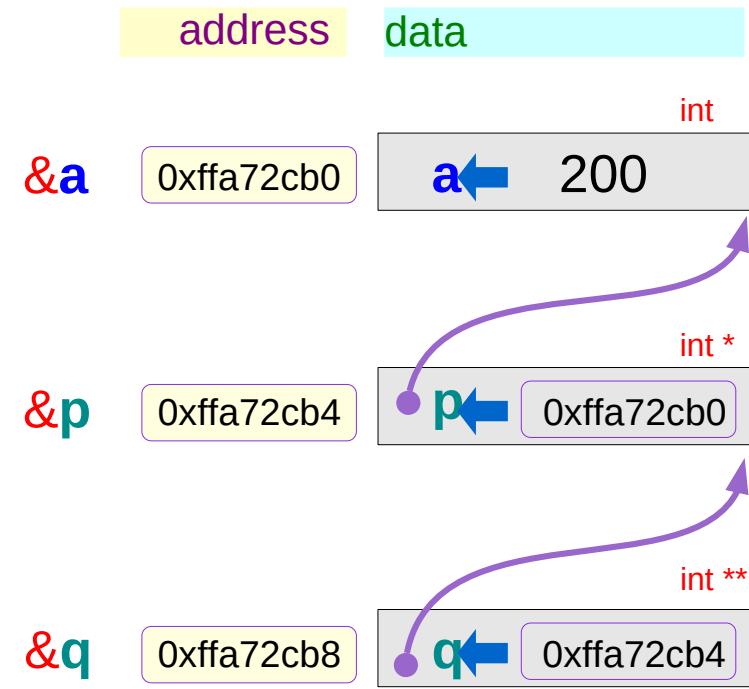
```
int    a = 200;  
int *  p = &a;
```

```
&p → 0ffa72cb4  
p → 0ffa72cb0  
*p → 200
```

# Pointer variable **q** of the type **int \*\***



```
int    a = 200;  
int * p = &a;  
int ** q = &p;
```



```
&q ← 0ffa72cb8  
q ← 0ffa72cb4  
*q ← 0ffa72cb0  
**q ← 200
```

# Pointer variable example codes

```
// t.c
#include <stdio.h>

int main(void) {
    int a = 200;
    int *p = &a;
    int **q = &p;

    printf("&a=%p a=%d \n", &a, a);
    printf("&p=%p p=%p \n", &p, p);
    printf("&q=%p q=%p \n", &q, q);
}
```

```
gcc -Wall -m32 t.c
./a.out
```

```
&a= 0ffa72cb0 a= 200
&p= 0ffa72cb4 p= 0ffa72cb0
&q= 0ffa72cb8 q= 0ffa72cb4
```

```
printf(" &p=%12p \n", &p);
printf(" p=%12p \n", p);
printf(" *p=%12d \n", *p);
printf("\n");

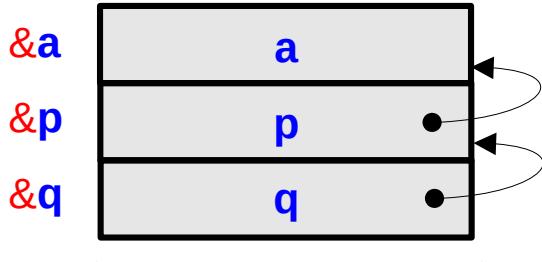
printf(" &q=%12p \n", &q);
printf(" q=%12p \n", q);
printf(" *q=%12p \n", *q);
printf(" **q=%12d \n", **q);
```

```
&p= 0ffa72cb4
p= 0ffa72cb0
*p= 200

&q= 0ffa72cb8
q= 0ffa72cb4
*q= 0ffa72cb0
**q= 200
```

# Byte addresses of variables

abstract  
address



`sizeof(a) = sizeof(int) =  
sizeof(p) = sizeof(int *) =  
sizeof(q) = sizeof(int **) =`  
4 bytes

```
int a = 200;  
int *p = &a;  
int **q = &p;
```

`&a= 0xffa72cb0 → value (&a)  
&p= 0xffa72cb4 → value (&p)  
&q= 0xffa72cb8 → value (&q)`

here, the values of **&a**, **&p**, **&q**  
are displayed

byte  
address

0xffa72cb0

0xffa72cb4

0xffa72cb8

MSByte

LSByte

00	00	00	c8
ff	a7	2c	b0
ff	a7	2c	b4

address value

**byte address** of  
the least significant byte  
in a little endian system

byte  
address

0xffa72cb0

0xffa72cb1

0xffa72cb2

0xffa72cb3

0xffa72cb4

0xffa72cb5

0xffa72cb6

0xffa72cb7

0xffa72cb8

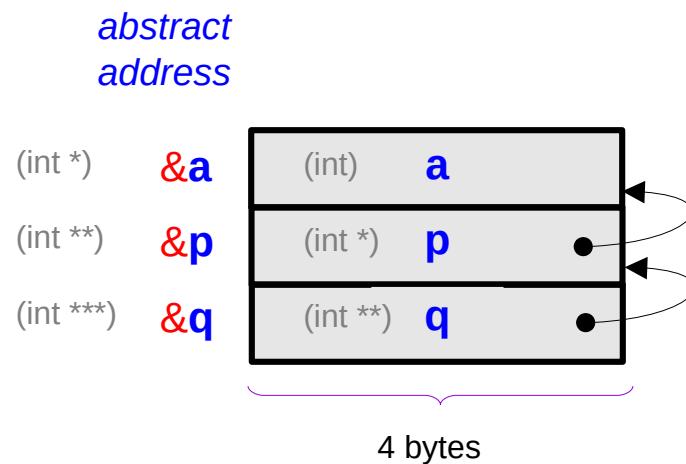
0xffa72cb9

0xffa72cba

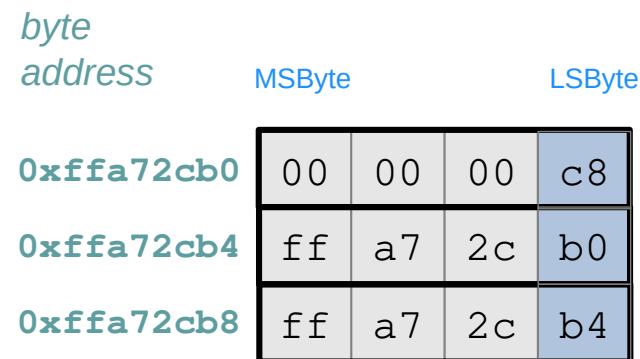
0xffa72ccb

c8
00
00
00
b0
2c
a7
ff
b4
2c
a7
ff

# Abstract vs. byte addresses



*abstract address contains information of **type**, **size**, **value***



*byte address contains **value** information only*

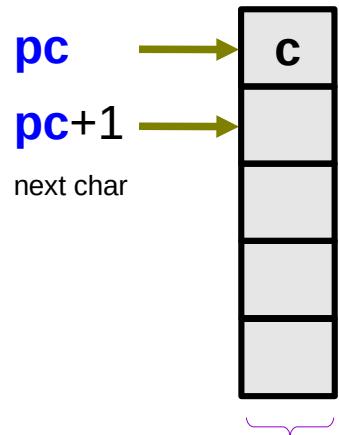
Abstract Address	Type	Size	Value (address value)
<b>&amp;a</b>	int *	4 bytes	<b>0xffa72cb0</b>
<b>&amp;p</b>	int **	4 bytes	<b>0xffa72cb4</b>
<b>&amp;q</b>	int ***	4 bytes	<b>0xffa72cb8</b>

Byte Address	Value
<b>value(&amp;a)</b>	<b>0xffa72cb0</b>
<b>value(&amp;p)</b>	<b>0xffa72cb4</b>
<b>value(&amp;q)</b>	<b>0xffa72cb8</b>

# Endianness and memory alignment

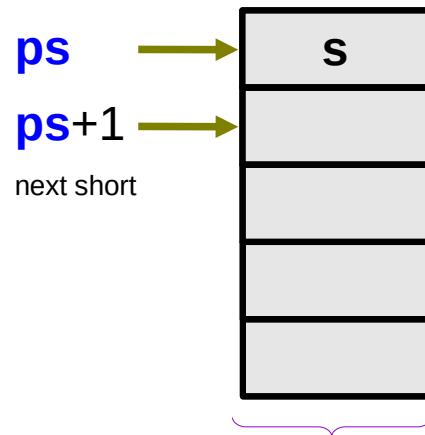
# Pointers and abstract addresses

char **c** ;  
char \***pc** = &**c** ;



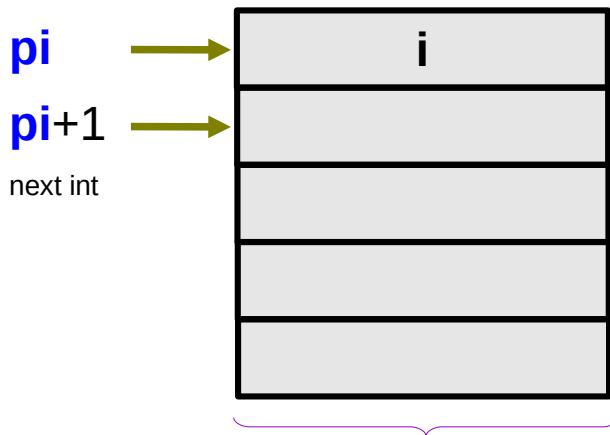
sizeof(\***pc**) =  
sizeof( **c** ) =  
sizeof( int ) = 1 byte

short **s** ;  
short \***ps** = &**s** ;



sizeof(\***ps**) =  
sizeof( **s** ) =  
sizeof( short ) = 2 bytes

int **i** ;  
int \***pi** = &**i** ;



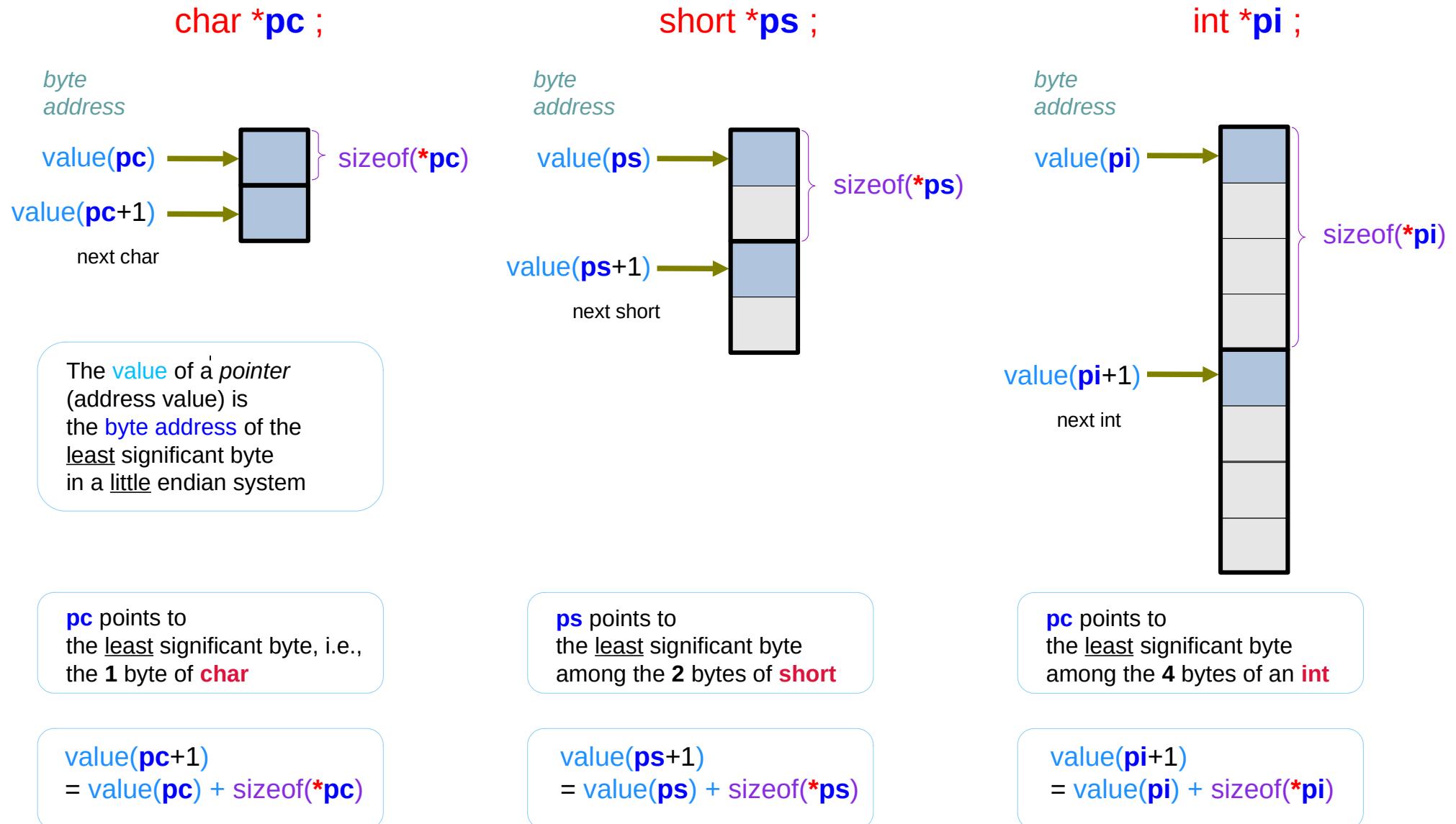
sizeof(\***pi**) =  
sizeof( **i** ) =  
sizeof( int ) = 4 bytes

sizeof(**pc**) = 4 or 8 bytes  
type(**pc**) = char \*  
value(**pc**) = byte address

sizeof(**ps**) = 4 or 8 bytes  
type(**ps**) = short \*  
value(**ps**) = byte address

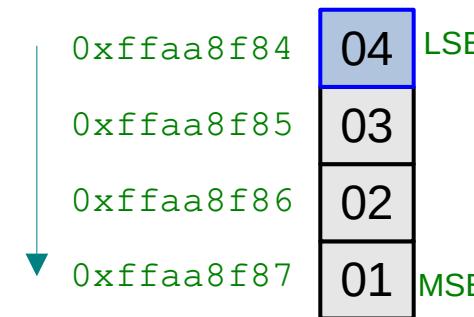
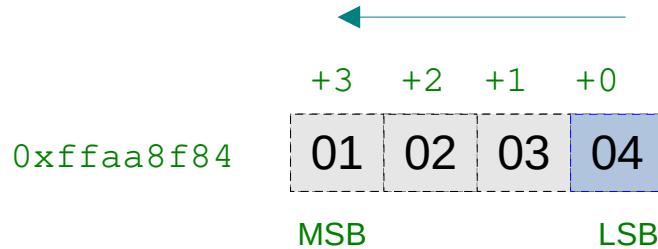
sizeof(**pi**) = 4 or 8 bytes  
type(**pi**) = int \*  
value(**pi**) = byte address

# Pointer values and byte addresses

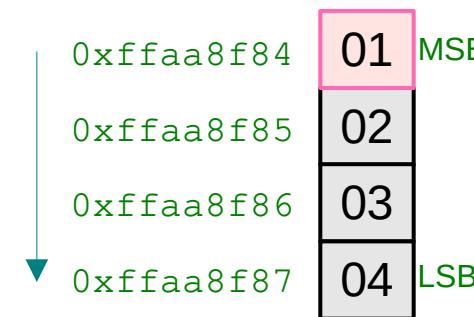
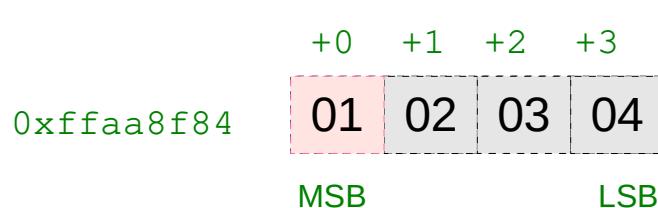


# Little Endian and Big Endian

MSB      LSB  
**int i = 0x01020304;**



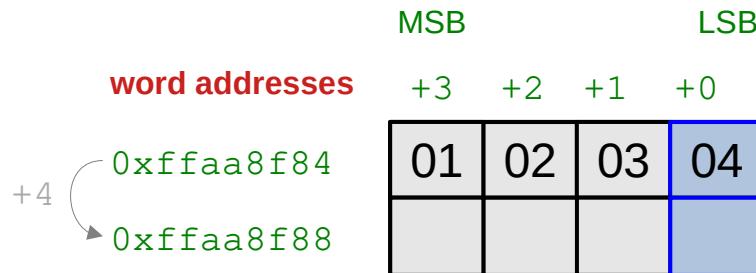
**Little Endian**  
stores LSByte first



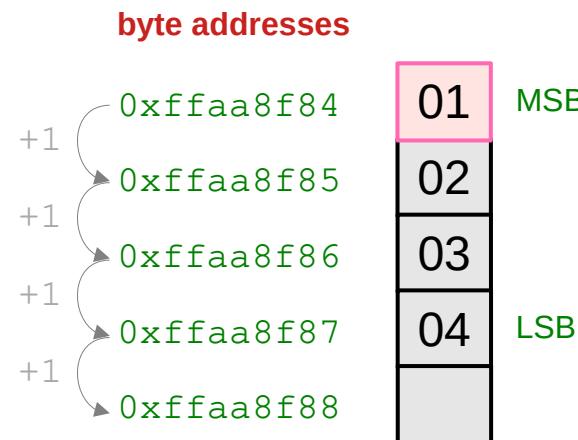
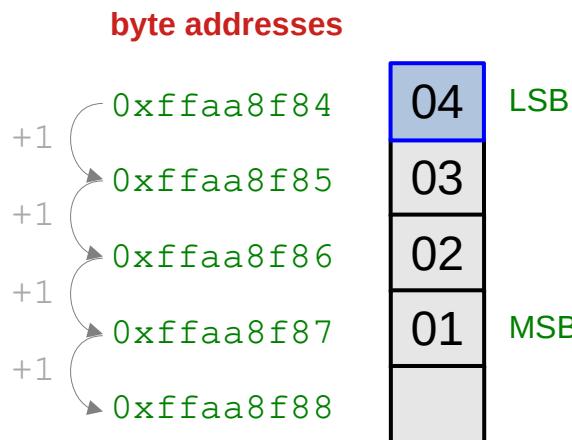
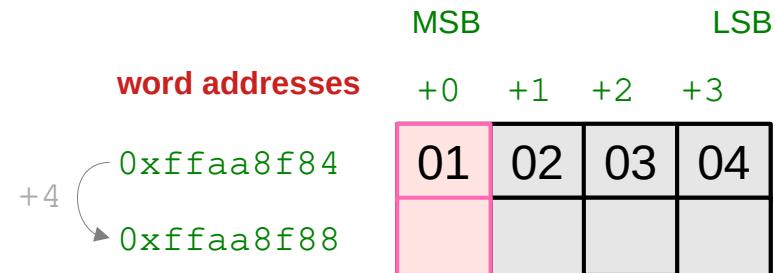
**Big Endian**  
stores MSByte first

# Word addresses and byte addresses

**LittleEndian** store LSByte first

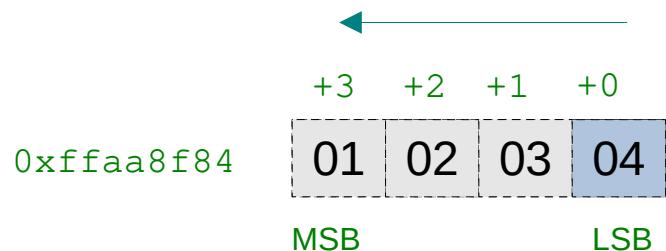


**BigEndian** store MSByte first



# Little Endian Examples

MSB                    LSB  
**int i = 0x01020304;**



**Little Endian**

store LSB first

Least Significant Byte (LSB)

```
MSB                    LSB  
int i     = 0x01020304;  
char *p = (char *)&i;  
  
printf("%p \n", &i);  
  
printf("%p : %x \n", p+0, p[0]);  
printf("%p : %x \n", p+1, p[1]);  
printf("%p : %x \n", p+2, p[2]);  
printf("%p : %x \n", p+3, p[3]);
```

0xffaa8f84		
LSB address	0xffaa8f84 : 4	LSB data
	0xffaa8f85 : 3	
	0xffaa8f86 : 2	
MSB address	0xffaa8f87 : 1	MSB data

assume 32-bit CPU  
32-bit address  
32-bit data (1 word, 4 bytes)

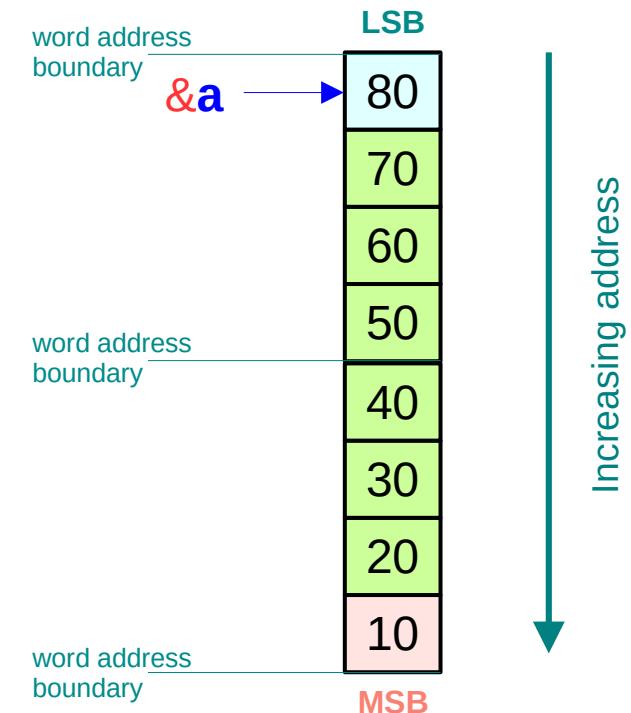
# Storing values in a little endian system

long      **a** = 0x1020304050607080 ;

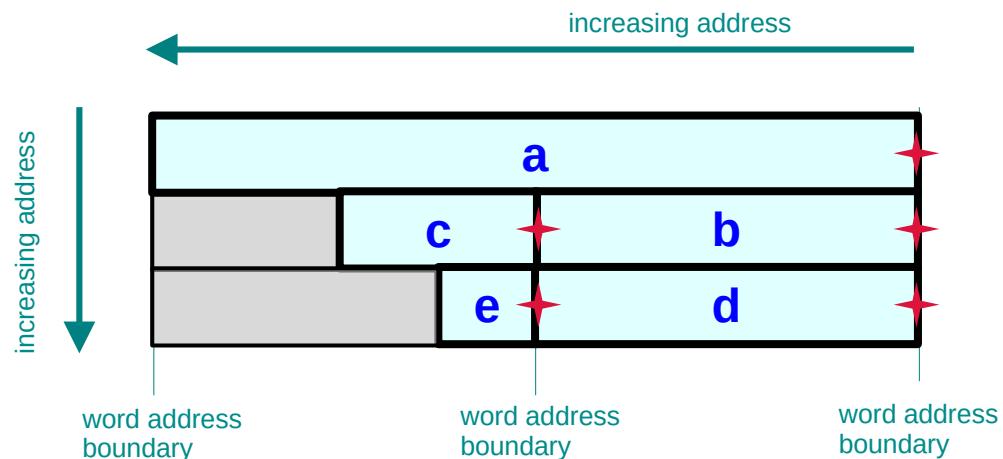
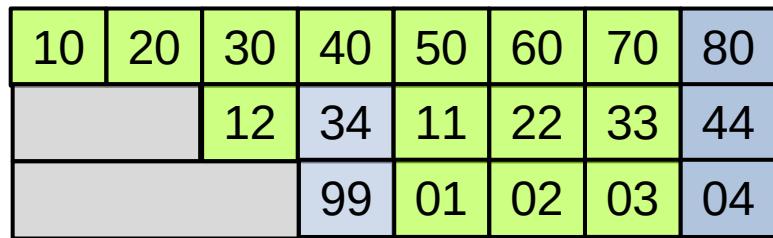
MSB (Most Significant Byte)  
LSB (Least Significant Byte)

## LittleEndian System

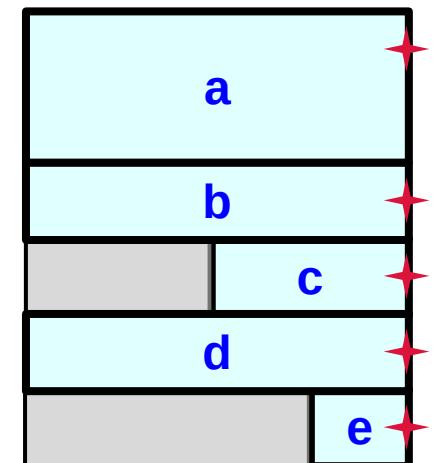
store the LSB first



# Memory alignment in a little endian system

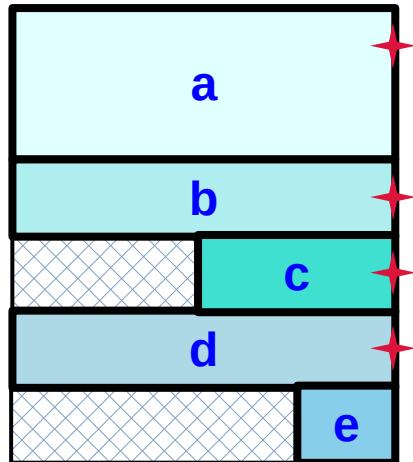


long      **a** = 0x1020304050607080 ;  
              MSB                            LSB  
  
int        **b** = 0x11223344 ;  
              MSB    LSB  
  
short      **c** = 0x1234 ;  
              MSB  
  
int        **d** = 0x01020304 ;  
              MSB                            LSB  
              LSB  
  
char      **e** = 0x99 ;

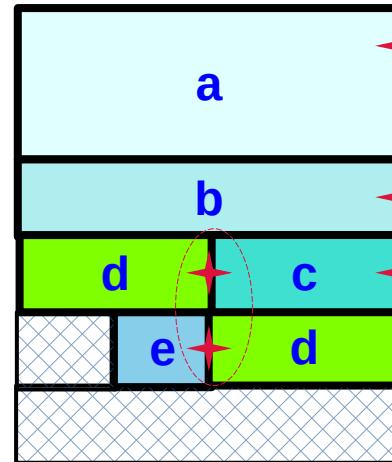


# Memory alignment – performance

## Memory alignment



## No memory alignment



saves memory  
loses performance :  
accessing **d** may require  
two memory cycles  
instead of one  
memory hardware constraints

long	<b>a</b> = 0x1020304050607080	MSB	LSB	;	
int	<b>b</b> =	0x11223344	MSB	LSB	;
short	<b>c</b> =	0x1234	MSB	LSB	;
int	<b>d</b> =	0x01020304	MSB	LSB	;
char	<b>e</b> =	0x99	MSB	LSB	;

# Pointer Type Cast

# Assigning pointer variables

long a ; &a address of a long value

int \* p ; address of an int value p ←~~←~~ &a address of a long value

short \* q ; address of a short value q ←~~←~~ &a address of a long value

char \* r ; address of a char value r ←~~←~~ &a address of a long value

# Type cast pointers

long a;

address of a long value &a

int \* p ;

address of an int value

p

int address conversion

(int \*) &a

short \* q ;

address of a short value

q

short address conversion

(short \*) &a

char \* r ;

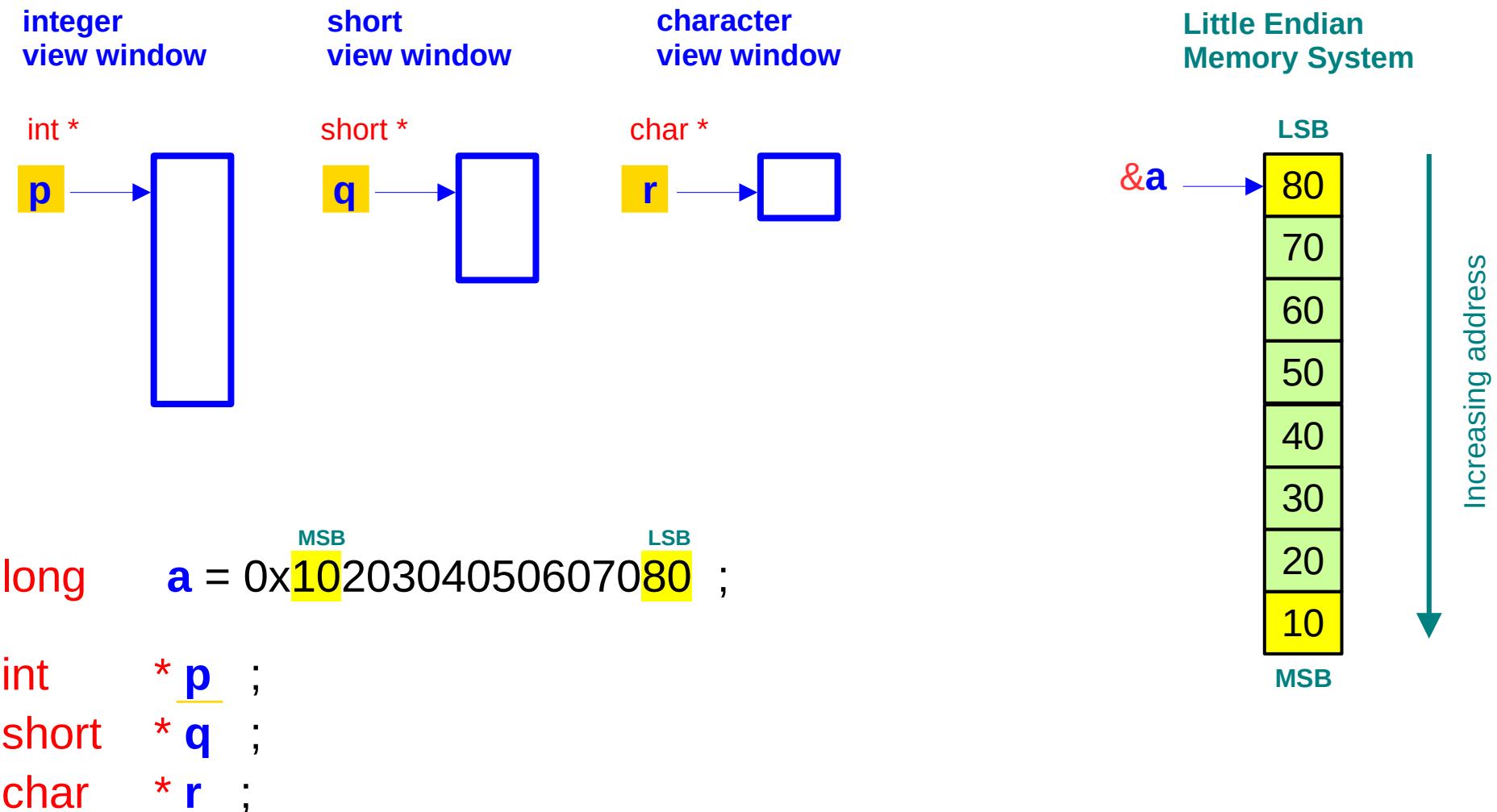
address of a char value

r

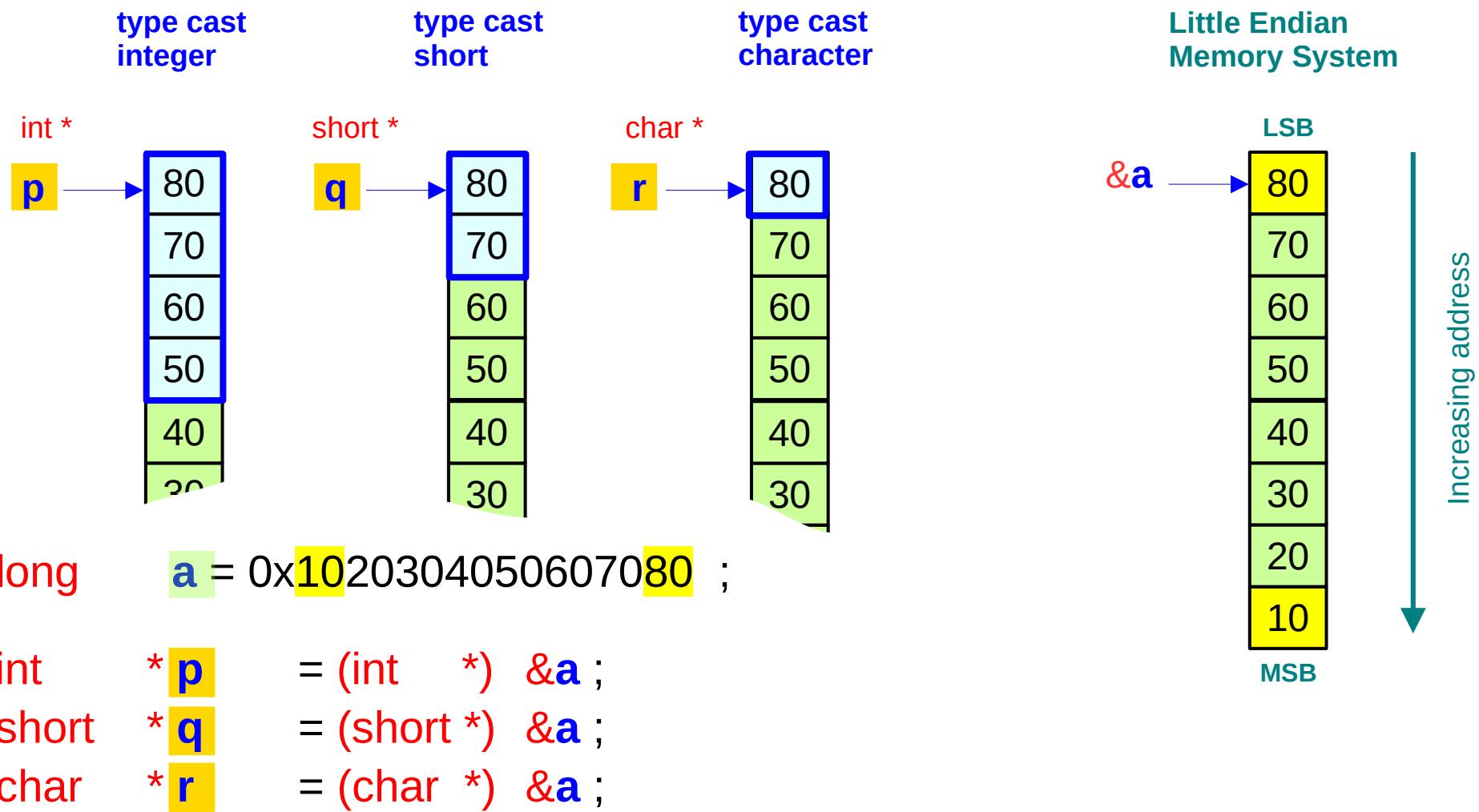
char address conversion

(char \*) &a

# View windows – type cast pointer

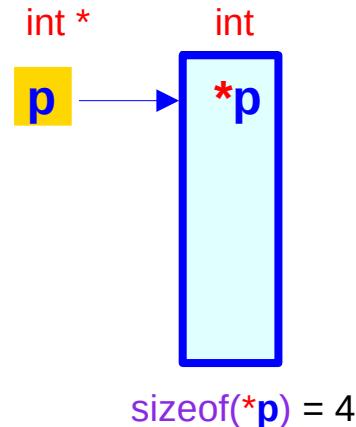


# Applying type cast pointers to a memory location

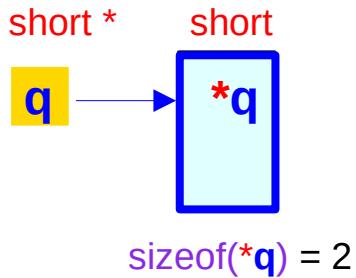


# De-referencing type cast pointers

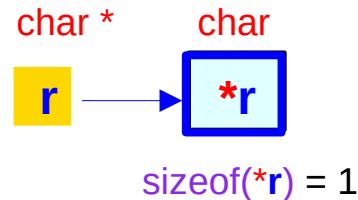
type cast  
integer pointer



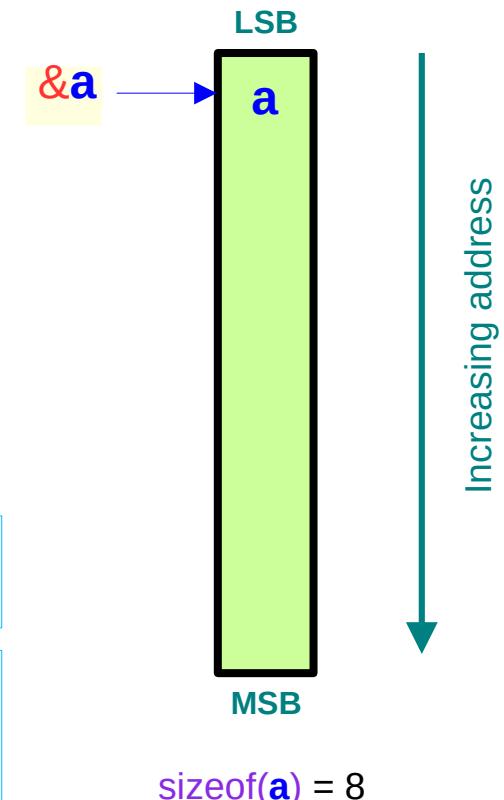
type cast  
short pointer



type cast  
character pointer



Little Endian  
Memory System

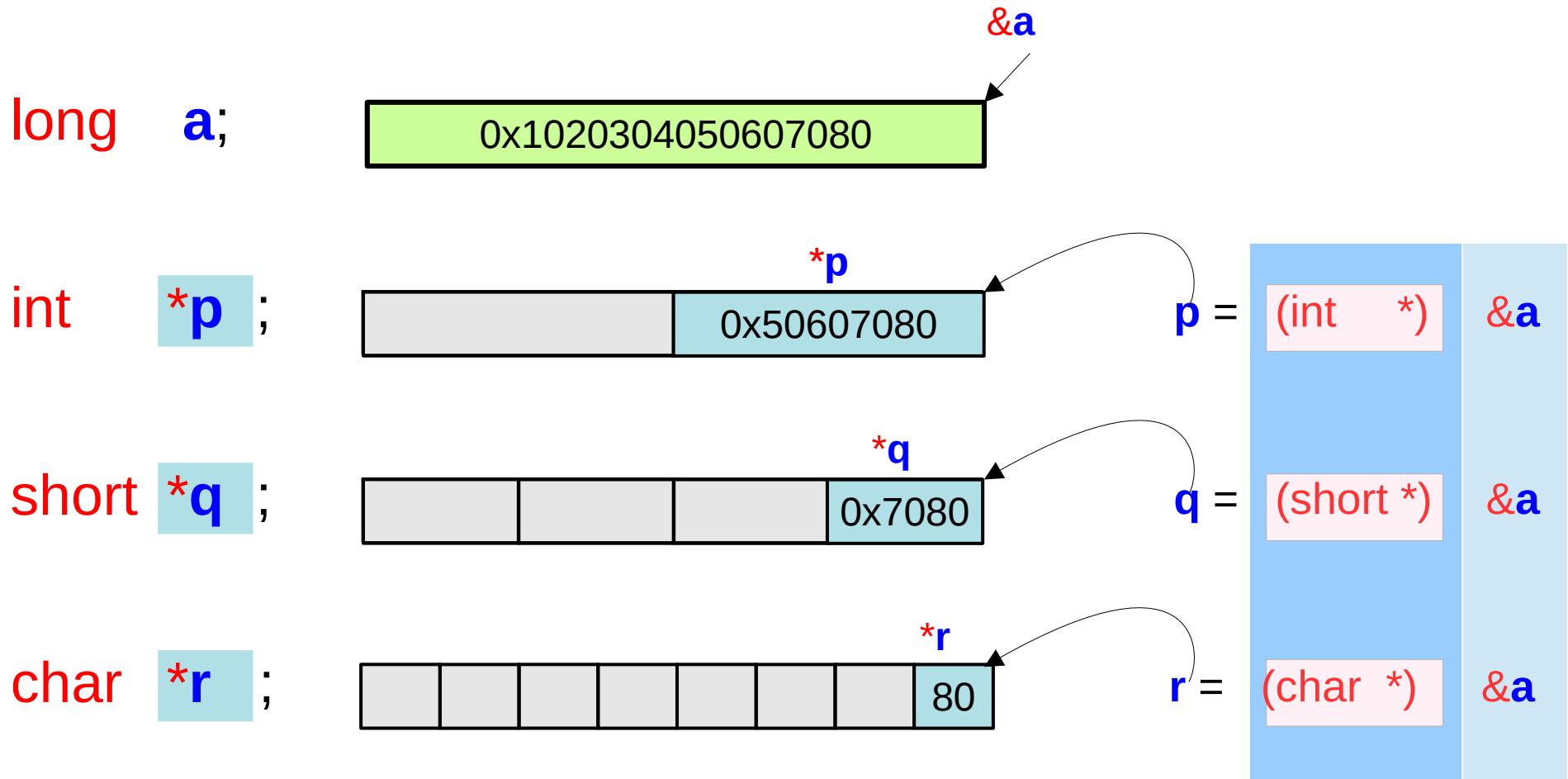


```
long a = 0x1020304050607080 ;
```

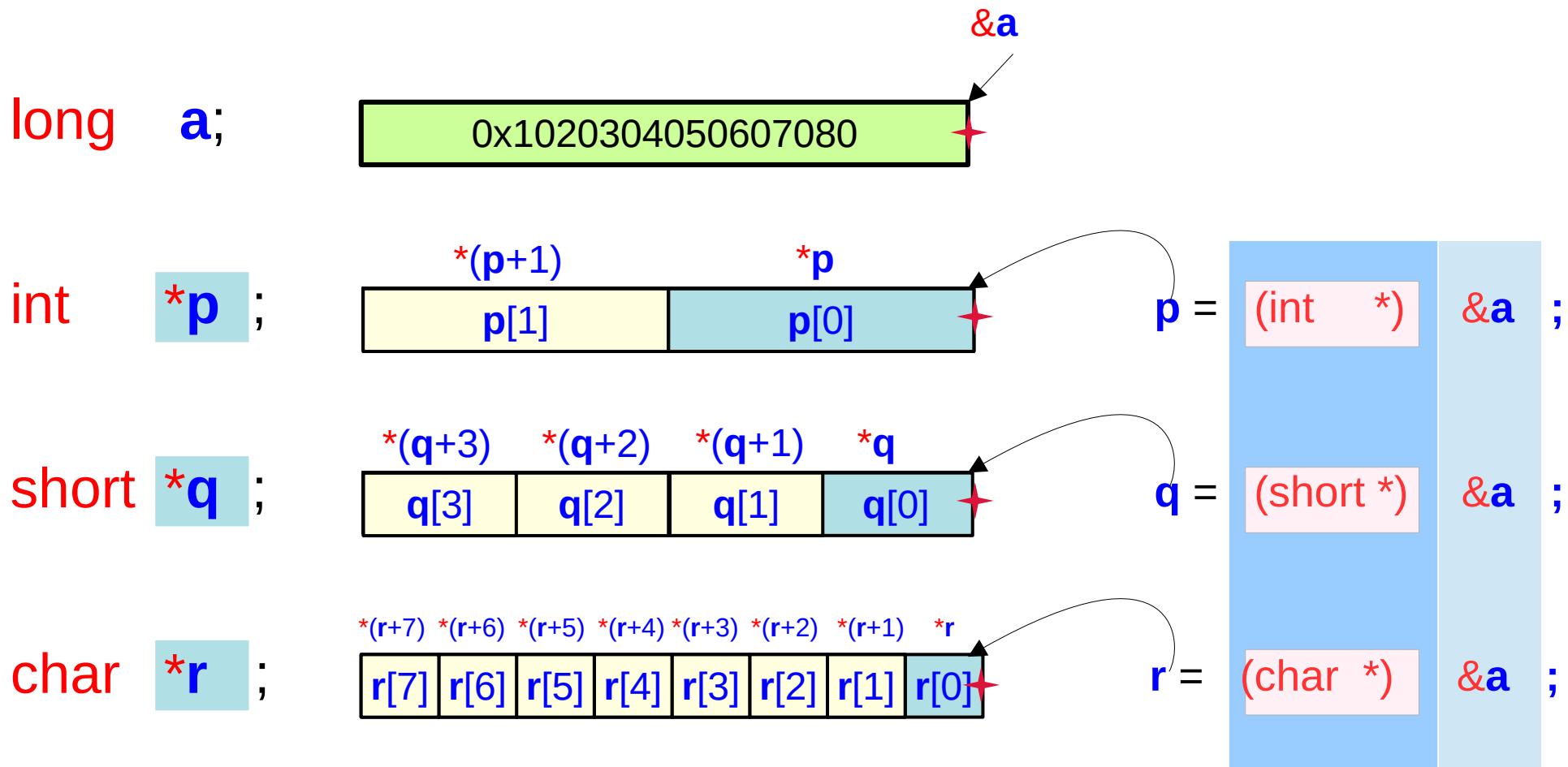
```
int *p = (int *) &a ;  
short *q = (short *) &a ;  
char *r = (char *) &a ;
```

```
*p ≡ 0x50607080  
*q ≡ 0x7080  
*r ≡ 0x80
```

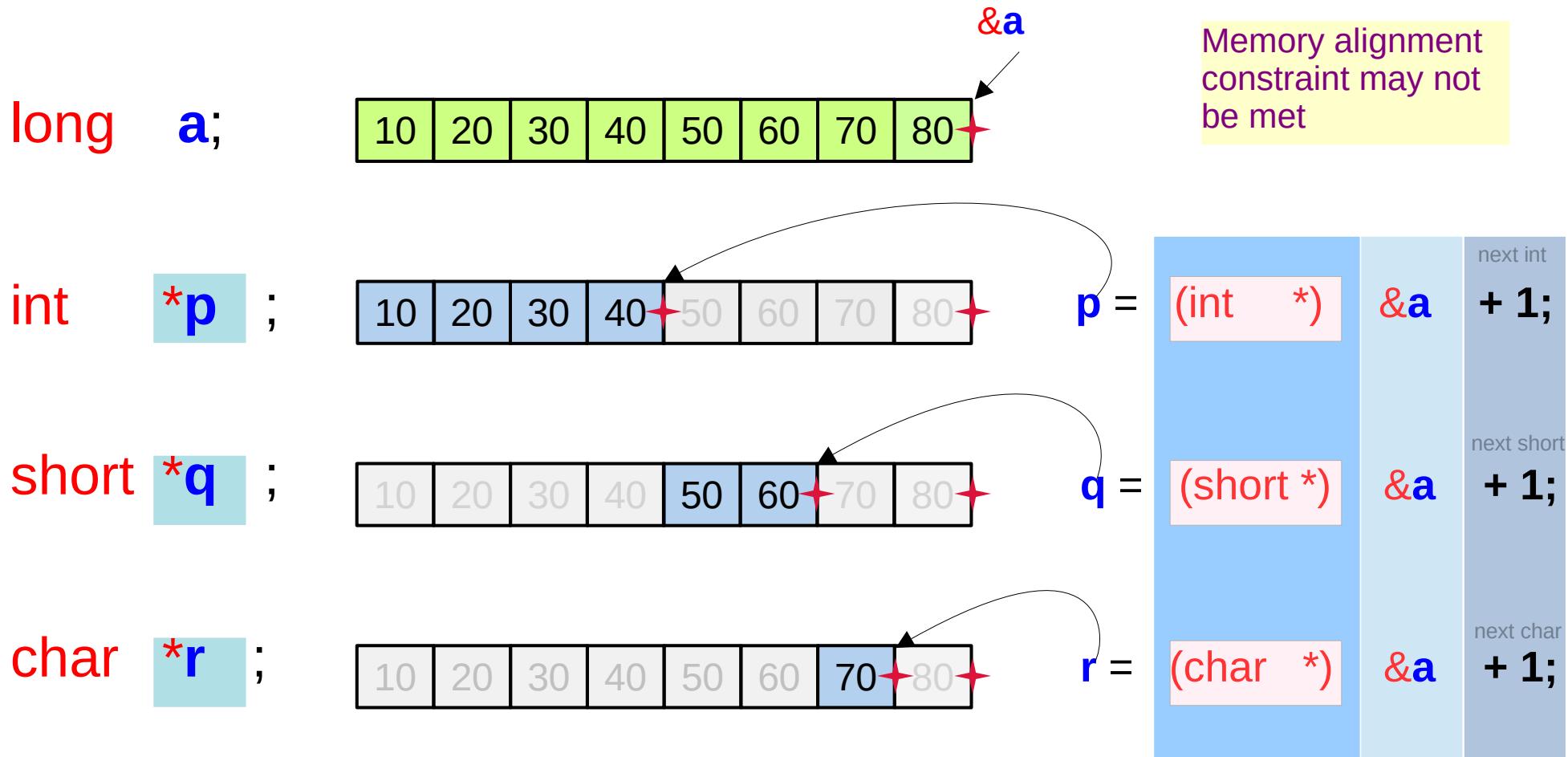
# Re-interpretation of memory data – case I



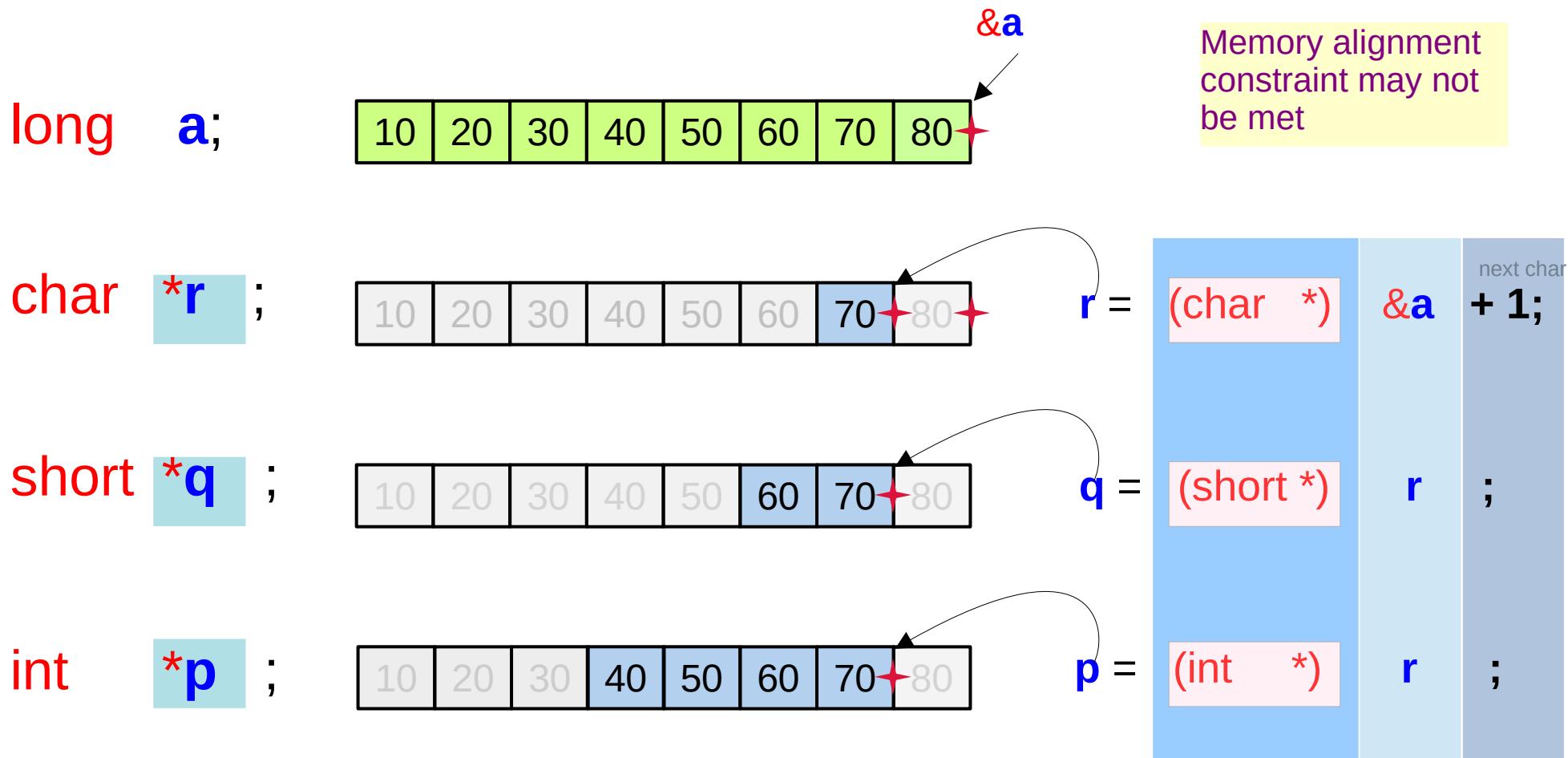
# Re-interpretation of memory data – case I



# Re-interpretation of memory data – case II



# Re-interpretation of memory data – case III



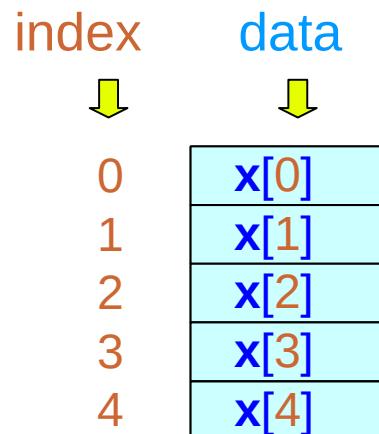
# Arrays

# Accessing array elements – using abstract addresses

```
int      x[5] ;
```

x holds the starting address of 5 consecutive int variables

5 int variables

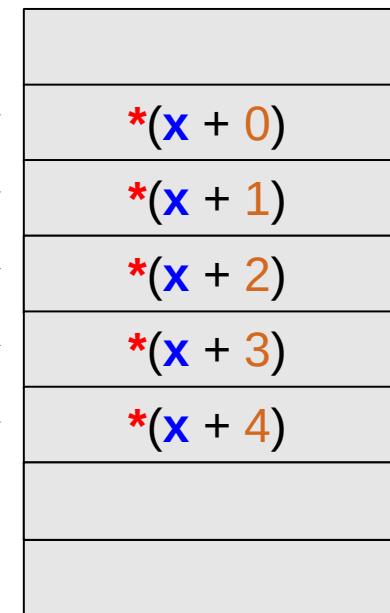


abstract  
address



data

x + 0 → \*(x + 0)  
x + 1 → \*(x + 1)  
x + 2 → \*(x + 2)  
x + 3 → \*(x + 3)  
x + 4 → \*(x + 4)



cannot change address x (constant)

# Accessing array elements – using byte addresses

int **x[5]** ;

**x** holds the *starting address* of 5 consecutive *int* variables

5 int variables

$$\begin{aligned} \text{value}(x + 0) &= \text{value}(x) + 0 * \text{sizeof}(*x) \\ \text{value}(x + 1) &= \text{value}(x) + 1 * \text{sizeof}(*x) \\ \text{value}(x + 2) &= \text{value}(x) + 2 * \text{sizeof}(*x) \end{aligned}$$

⋮

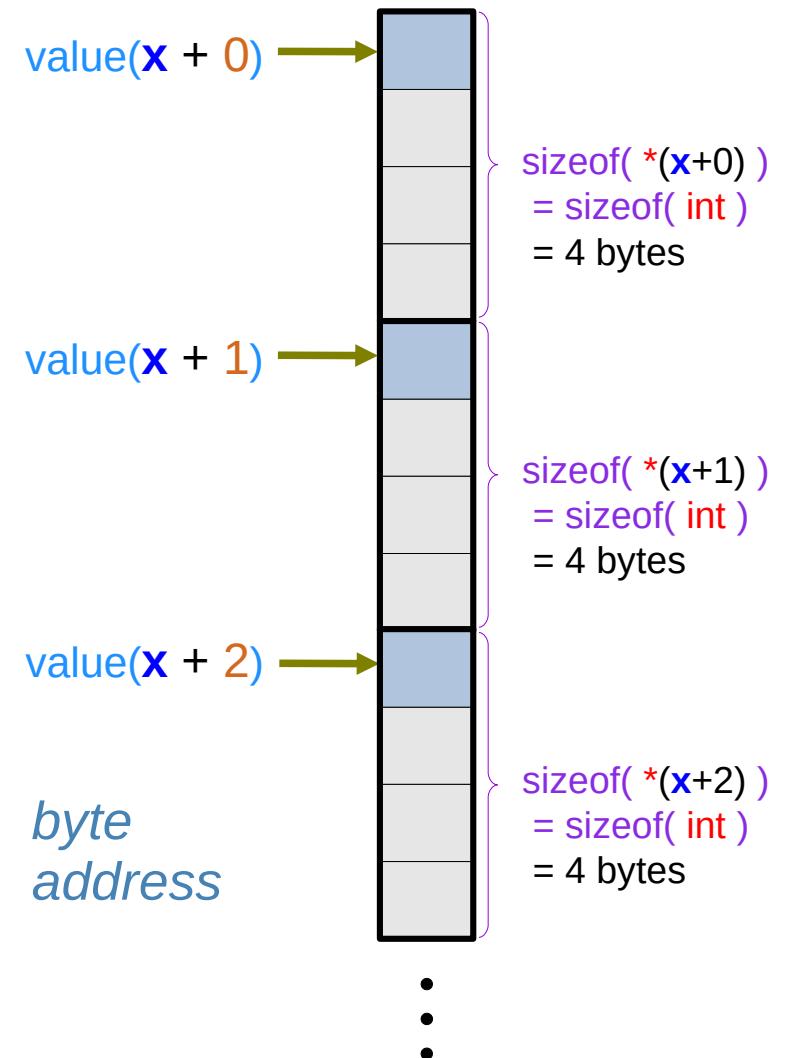
⋮

⋮

byte address

byte address

$\text{sizeof}(\text{int}) = 4$  bytes



## Two aspects of x

```
int      x[5];
```

x is an array of 5 integer elements

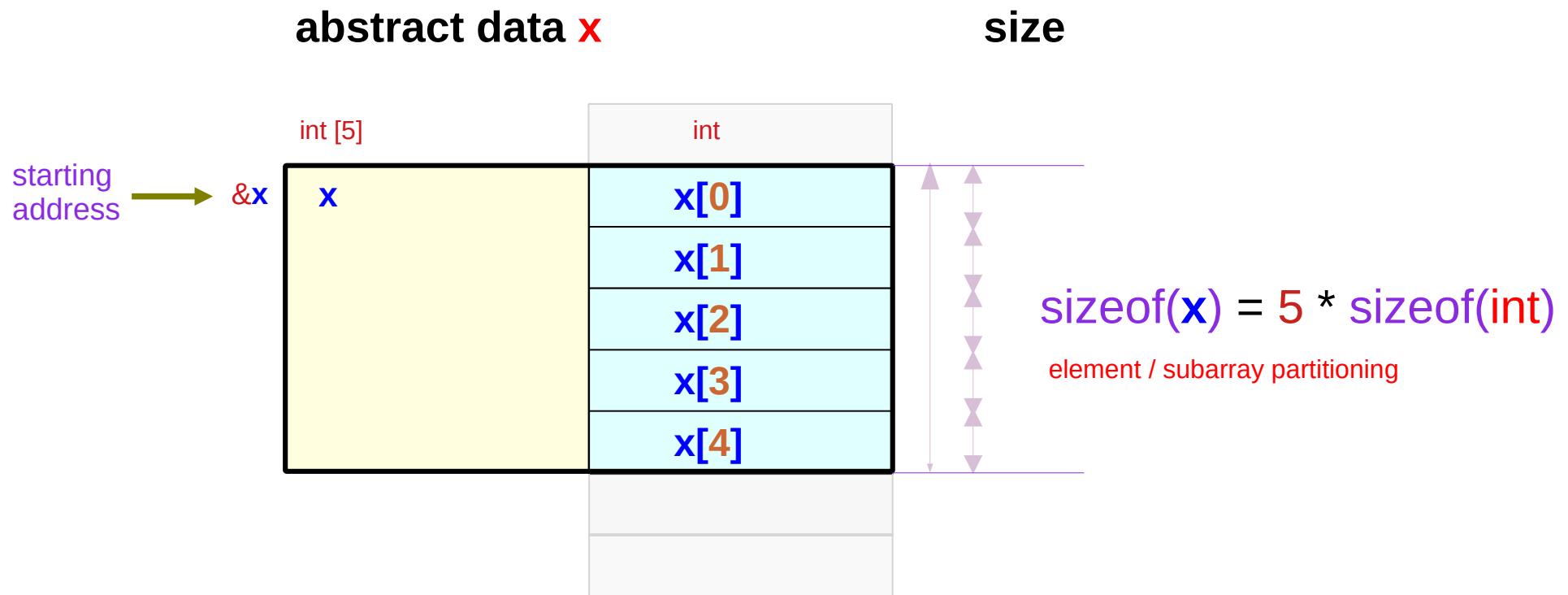
**abstract data x ..... size**

**pointer x.....starting address**

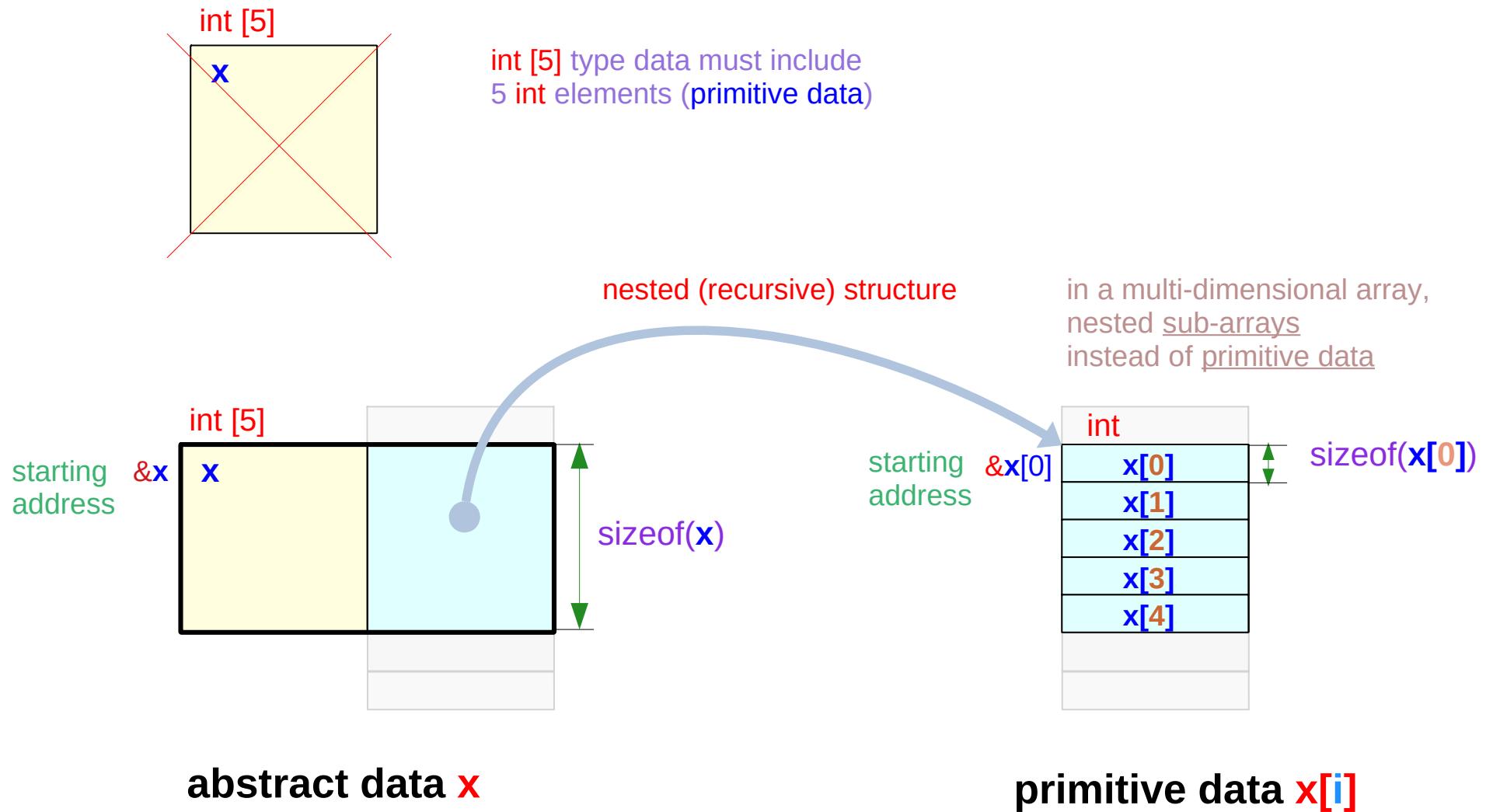
# Abstract data x

int **x[5];**

the size of x is the total size  
of 5 consecutive int variables



# Recursive structure



# Pointer x

int **x[5];**

**x** holds the starting address of 5 consecutive int variables

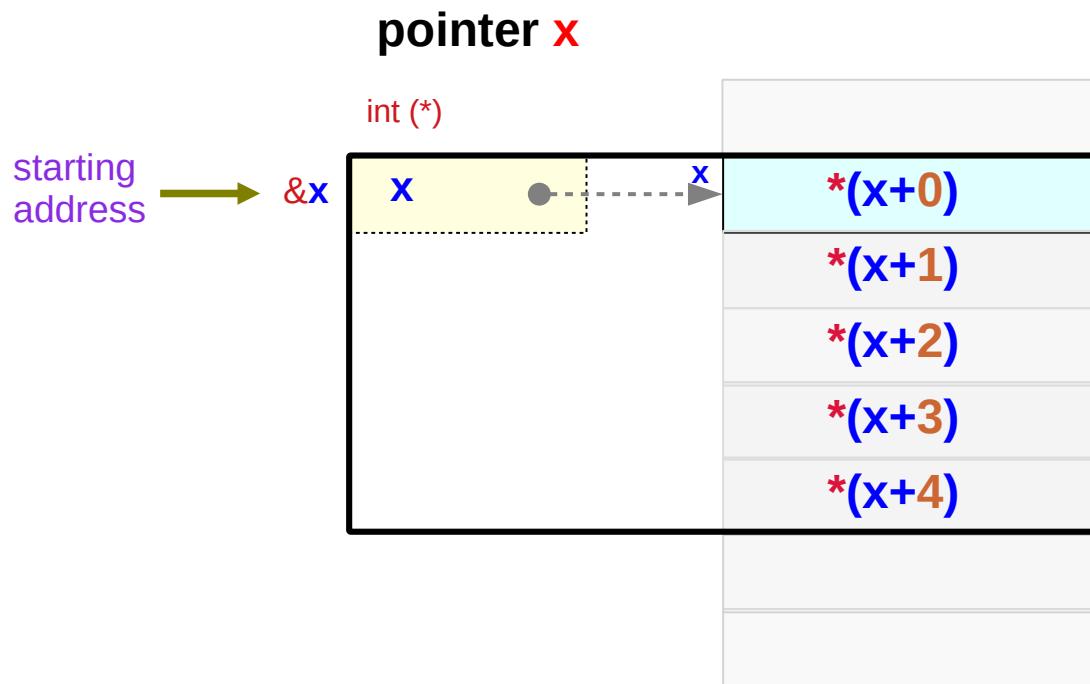
**x[i]  $\equiv$  \*(x+i)**

only for  $i = 0, \dots, 4$

**x[0]  $\equiv$  \*x**

**&x[0]  $\equiv$  x**

equivalence relations



**starting address**

the starting address of 5 consecutive int variables

**value(x) = value(&x[0])**

the starting address (**&x**) of the array **x**

**value(&x) = value(x)**

address replication

# Aggregate Data Type – 1-d Array

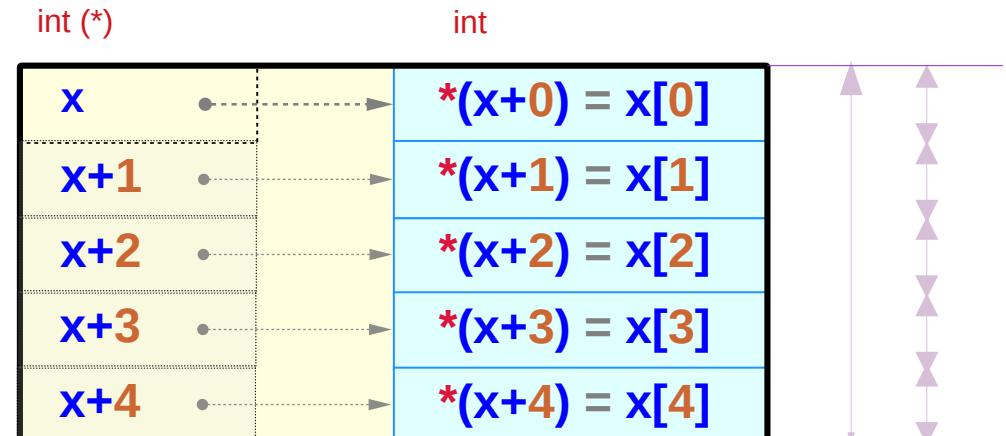
int      **x[5];**

**x** holds the starting address  
of 5 consecutive int variables

*the start address of each element :*

$$\begin{aligned} \text{value}(x+0) &= \text{value}(x) + 0 * \text{sizeof}(*x) \\ \text{value}(x+1) &= \text{value}(x) + 1 * \text{sizeof}(*x) \\ \text{value}(x+2) &= \text{value}(x) + 2 * \text{sizeof}(*x) \\ \text{value}(x+3) &= \text{value}(x) + 3 * \text{sizeof}(*x) \\ \text{value}(x+4) &= \text{value}(x) + 4 * \text{sizeof}(*x) \end{aligned}$$

**Aggregate data with 5 elements**



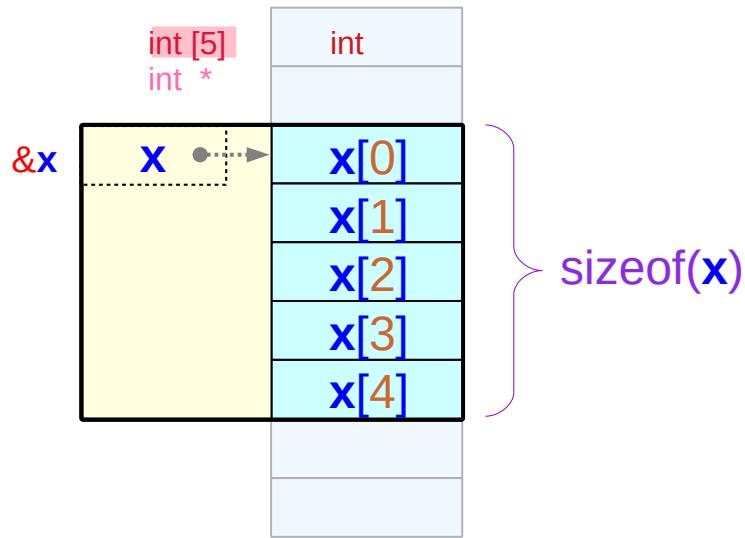
**value(x) = the start address of the array**

**sizeof(x) = sizeof(\*x) \* 5**

- An array **x** and a pointer **p**

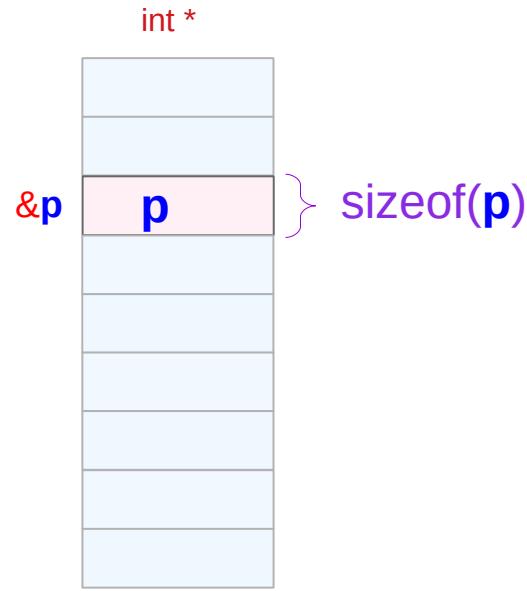
# Sizes of an array **x** and a pointer **p**

**int x [5] ;**



$$\text{sizeof}(x) = 5 * \text{sizeof}(\text{int})$$

**int \* p ;**



$$\begin{aligned}\text{sizeof}(p) &= \text{size of a pointer} \\ &= 4 \text{ bytes (32-bit system)} \\ &8 \text{ bytes (64-bit system)}\end{aligned}$$

# Address values of an array **x** and a pointer **p**

**int x [5] ;**

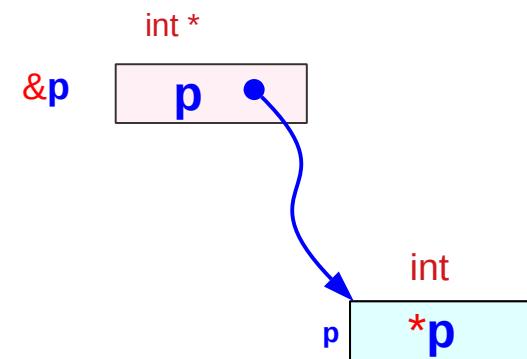
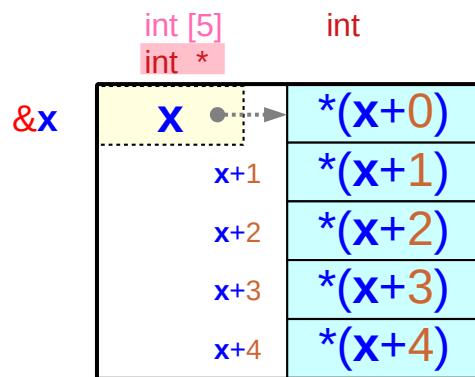
**x** : an array variable name (constant)

**value(x)** : the starting address of  
5 consecutive **int** variables

**int \* p ;**

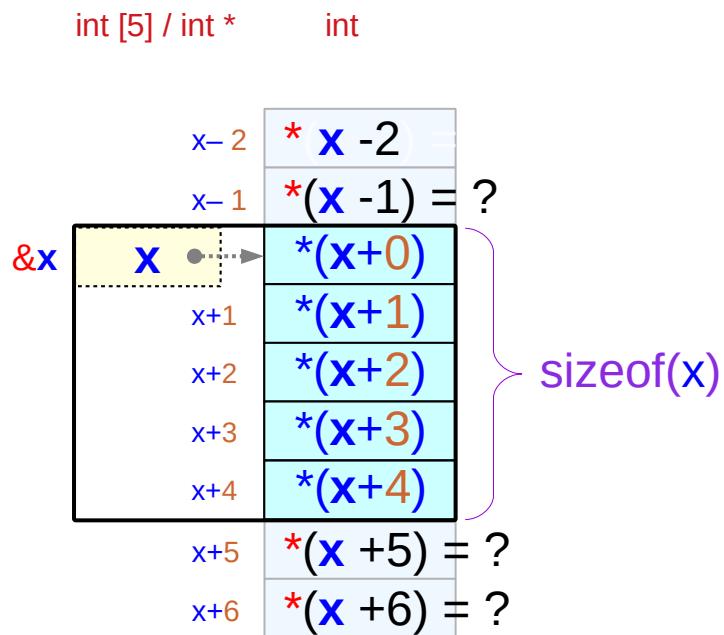
**p** : an variable name

**value(p)** : the address  
of an **int** variable

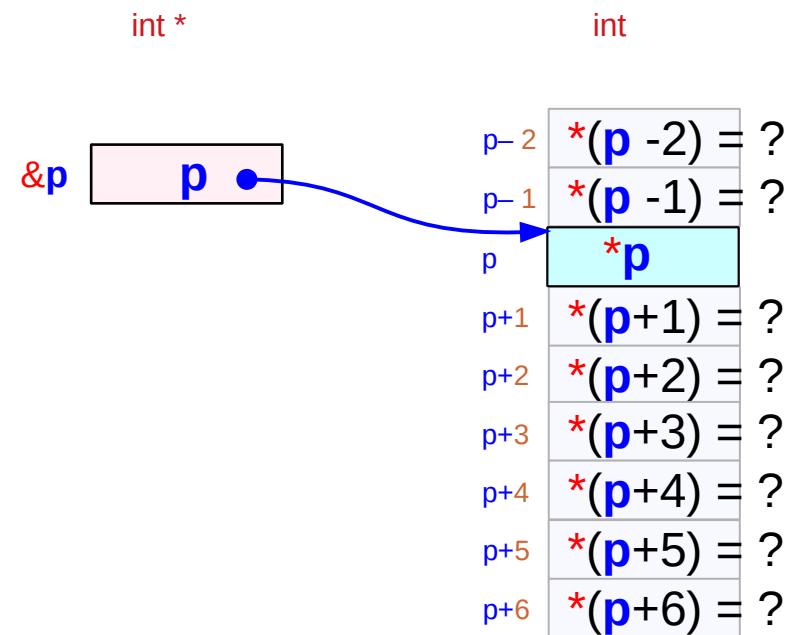


# Out of range index

int **x** [5] ;



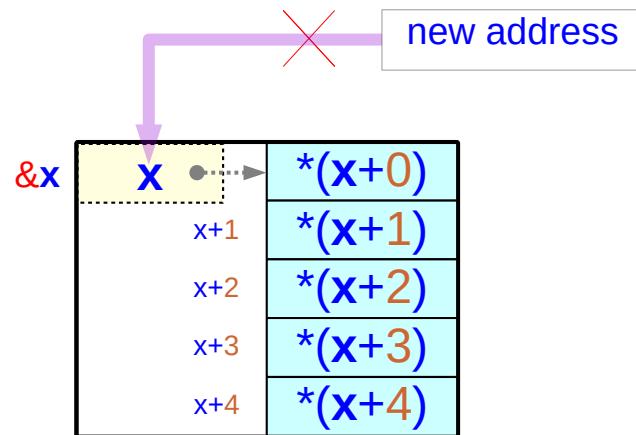
int \* **p** ;



A programmer's responsibility

# Assignment of a new address

`int x [5] ;`

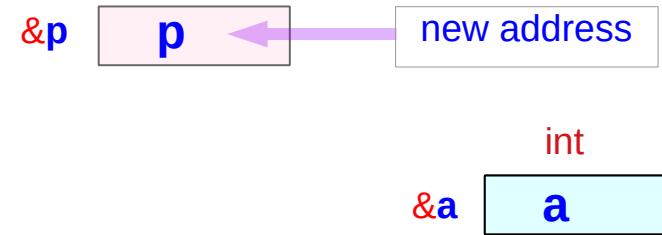


`x` is a constant variable  
(not an *lvalue*)

`x` and `&x` give the same  
value of `&x[0]` (address replication)

this address value(`x`) is assigned by  
the compiler and cannot be changed

`int * p ;`



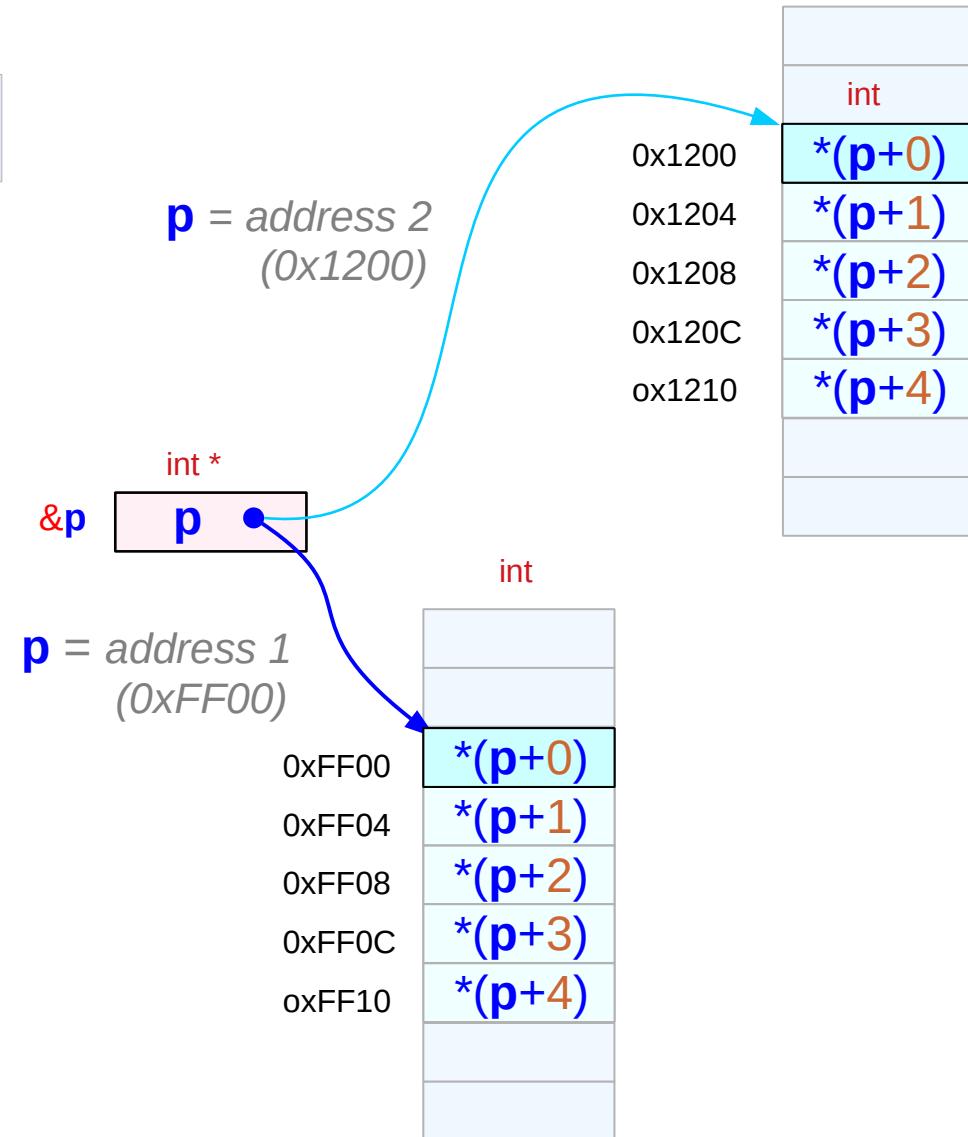
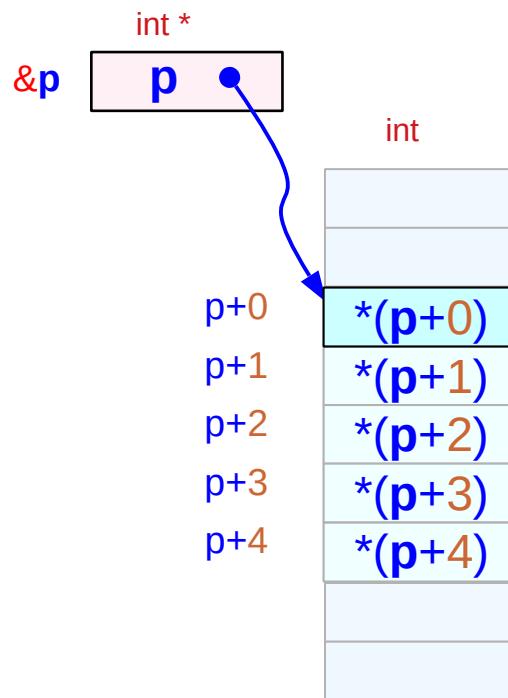
`p` is a pointer variable  
(a *lvalue*)

a memory location is allocated  
for a variable `p`

`value(p)` can be changed by an  
assignment statement, i.e,  
`p = &a`

# Pointer variable can point different locations

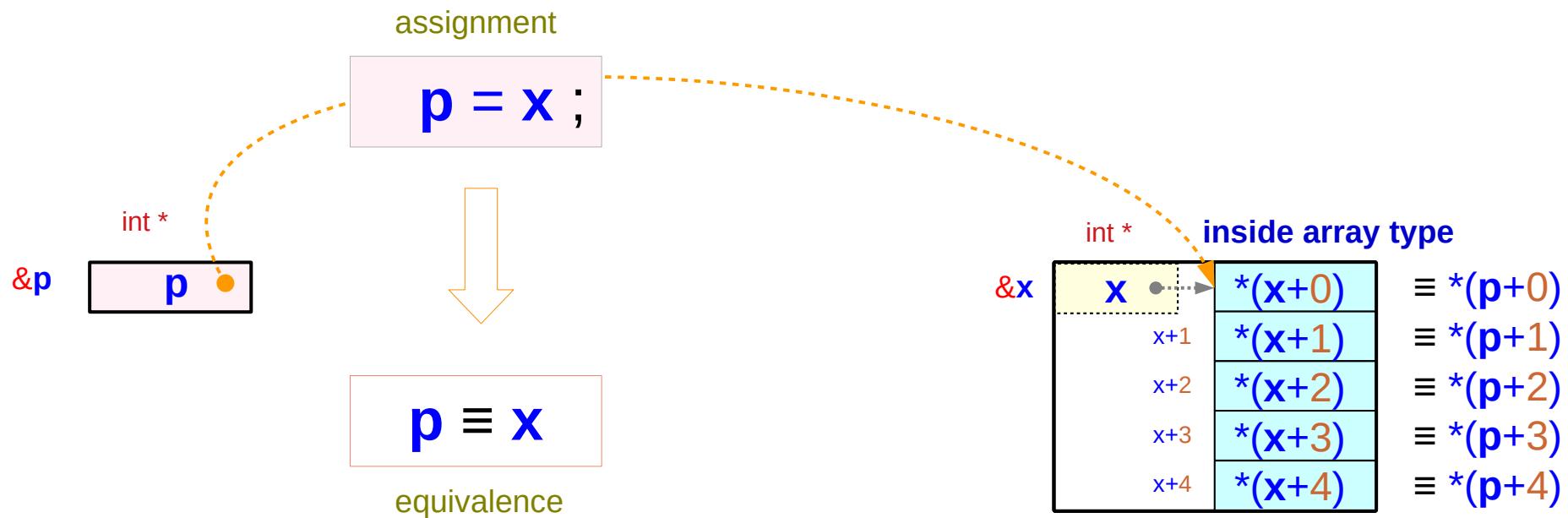
```
int * p ;
```



# Pointer to an array element

```
int (*p) ;
```

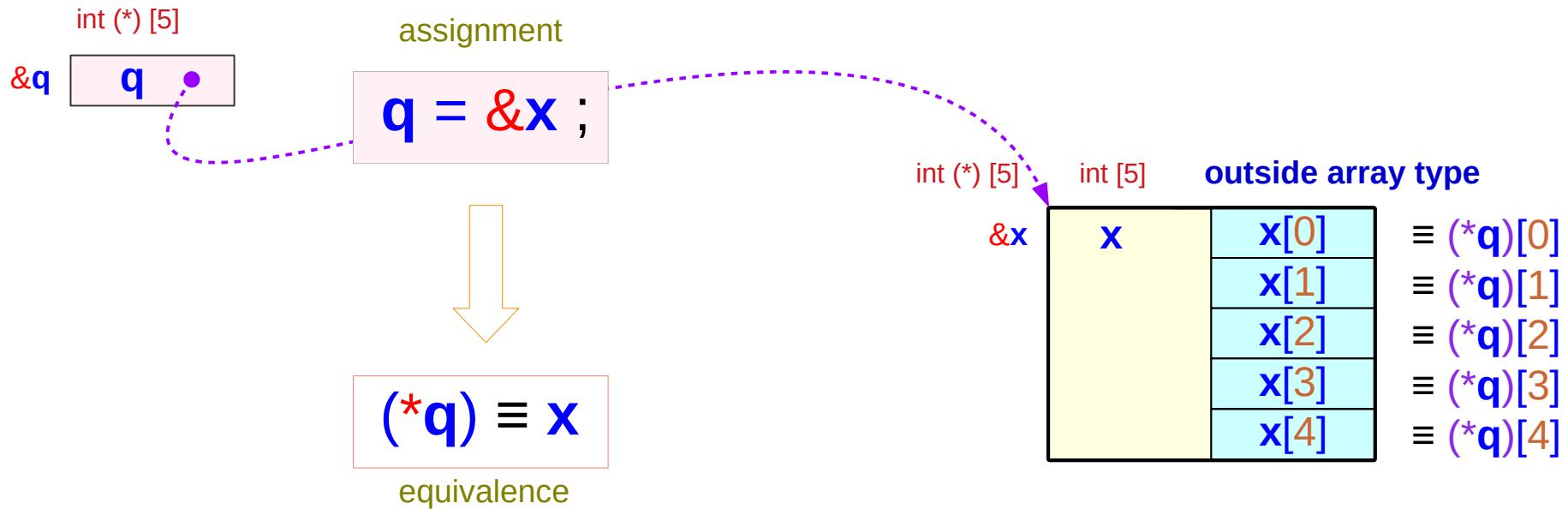
```
int x [5] ;
```



# Pointer to an array

```
int (*q) [5] ;
```

```
int x [5] ;
```



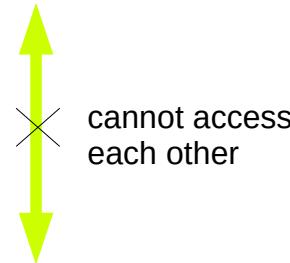
`q` is a pointer variable that  
points to an array with 5  
integer variables

# Pass by Reference

# Variable Scopes

```
int func1 (int a, int b)
{
    int i, int j;
    ...
    ...
}
```

i and j's  
variable scope



```
int main ()
{
    int x, int y;
    ...
    ...
    func1 ( 10, 20 );
    ...
    ...
}
```

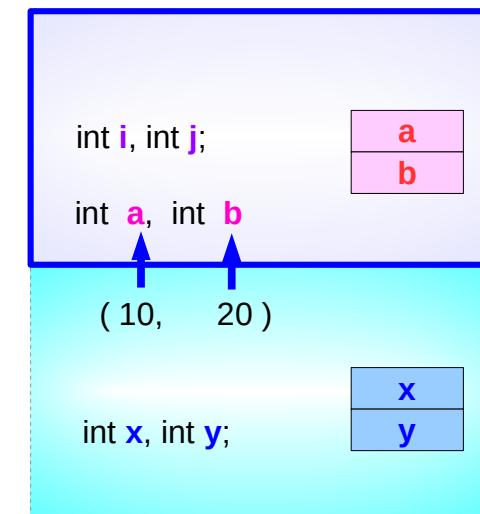
x and y's  
variable scope

Only **top** stack frame is active  
and its variable can be accessed

Communications are performed  
only through the **parameter** variables

func1's  
Stack  
Frame

main's  
Stack  
Frame



# Pass by Reference

```
int func1 (int* a, int* b)
{
    int i, int j;
    ...
    ...
    ...
}
```

x and y are made known to func1  
func1 can read / write x and y  
through their addresses

```
int main ()
{
    int x, int y;
    ...
    ...
    func1 ( &x, &y );
    ...
}
```

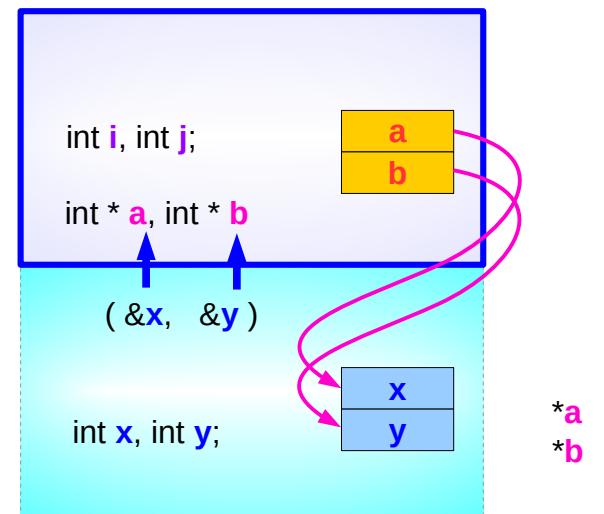
a = &x  
b = &y

x and y's  
variable scope

func1's  
Stack  
Frame

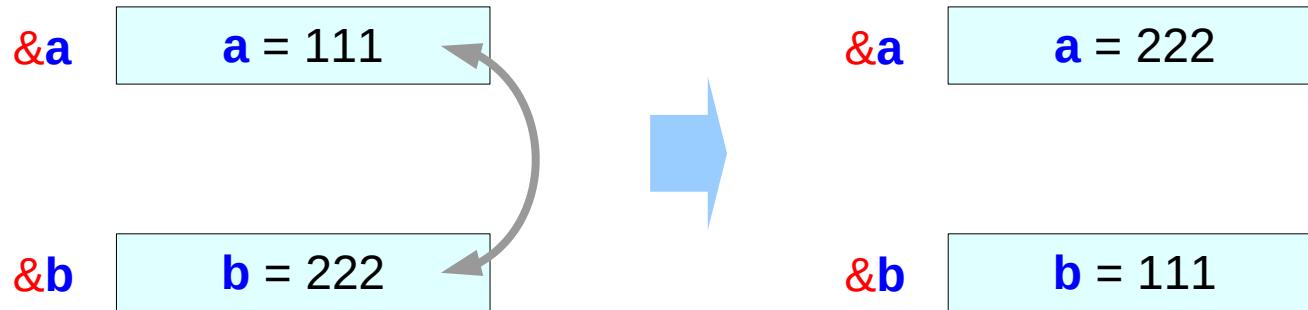
main's  
Stack  
Frame

\*a  
\*b



Example :  
swapping integers  
swapping pointers

# Swapping integers



```
int a, b;
```

```
swap( &a, &b );
```

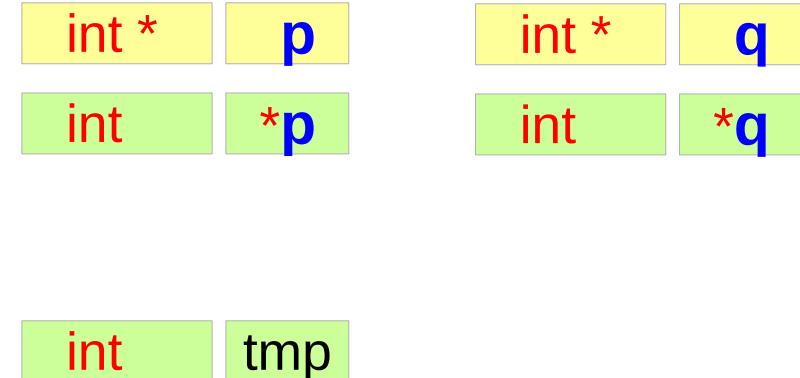
```
swap( int *, int * );
```

function call

function prototype

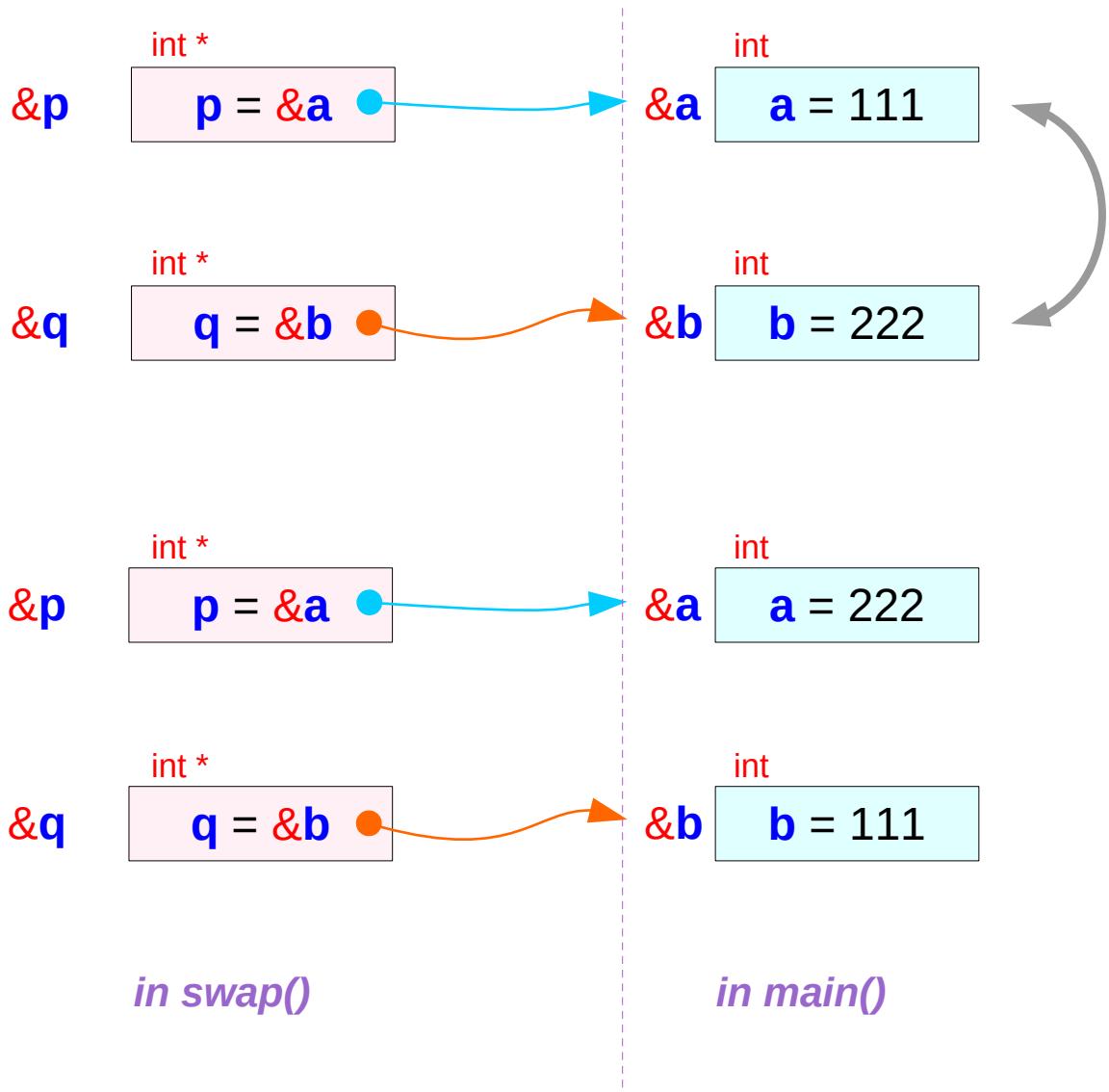
# Pass by integer reference

```
void swap(int *p, int *q) {  
    int tmp;  
  
    tmp = *p;  
    *p = *q;  
    *q = tmp;  
}
```



```
int a = 111, b = 222;  
...  
swap( &a, &b );
```

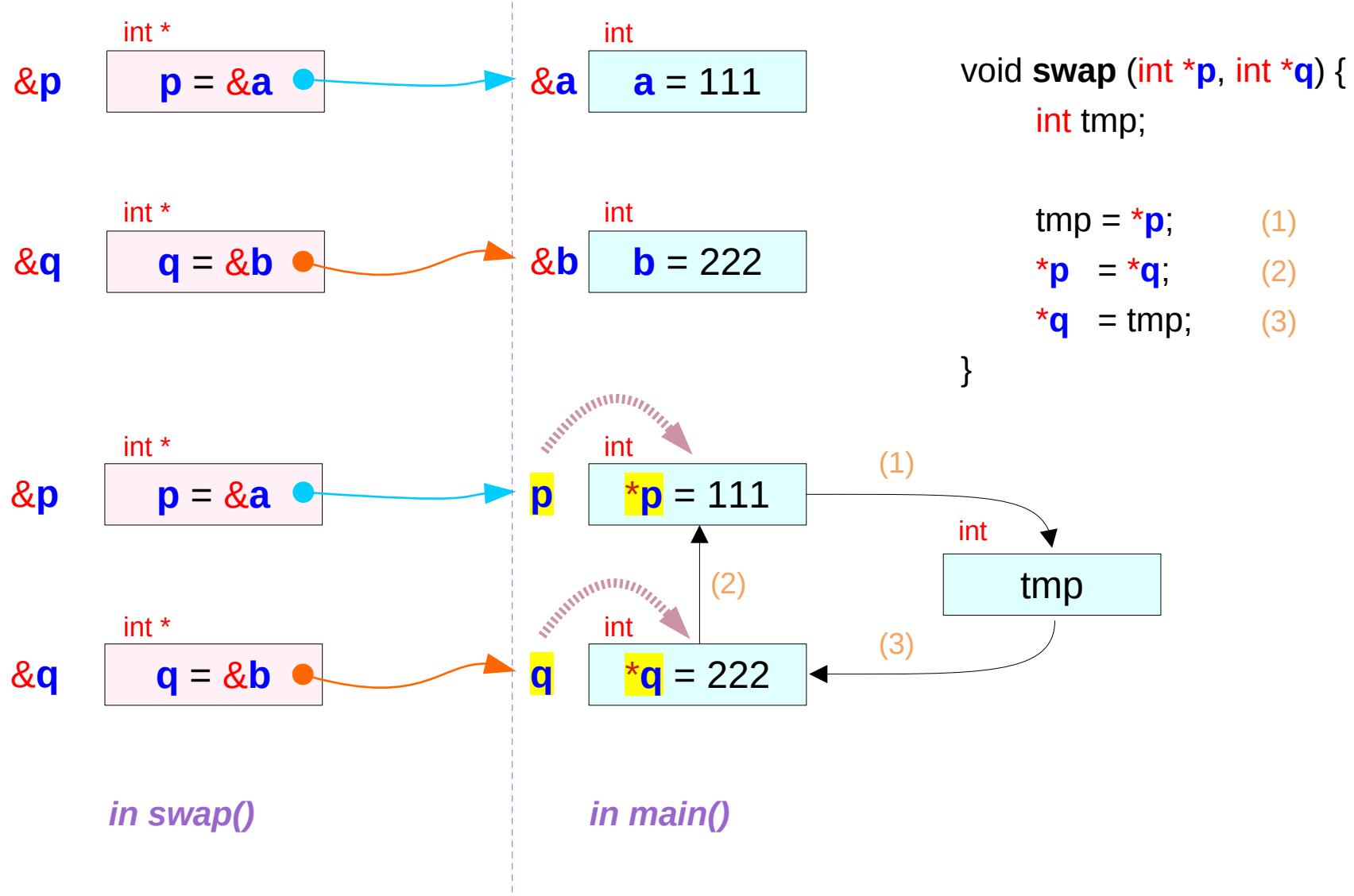
# Swapping integers



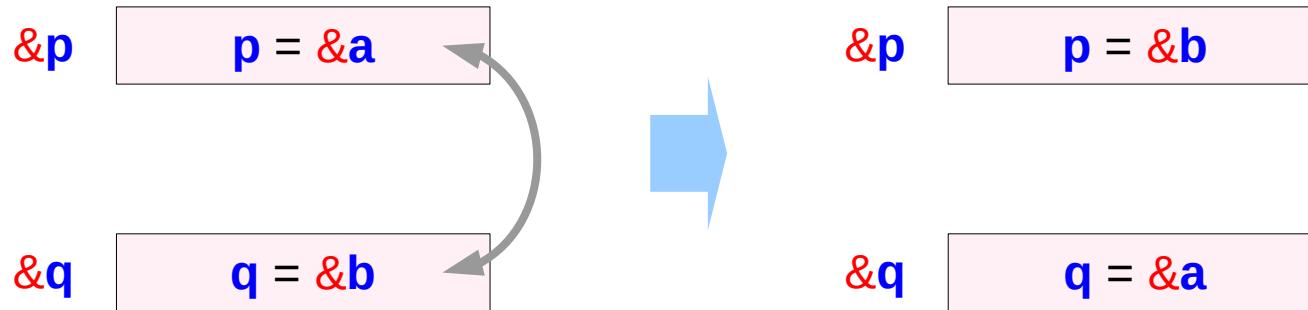
before

after

# Swapping integers via integer pointers



# Swapping integer pointers



The diagram illustrates the flow of pointers from a function call to its prototype. It consists of three horizontal rows of code:

- Top Row:** `int *p, *q ;`
- Middle Row:** `pswap ( &p, &q );` This row shows the actual arguments being passed to the function. The addresses of `p` and `q` are passed as arguments. Yellow arrows point from the `*p` and `*q` variables in the first row down to the `&p` and `&q` parameters in this row.
- Bottom Row:** `void pswap( int **, int ** );` This row shows the function prototype. It defines the function `pswap` to take two double pointers (`int **`) as parameters.

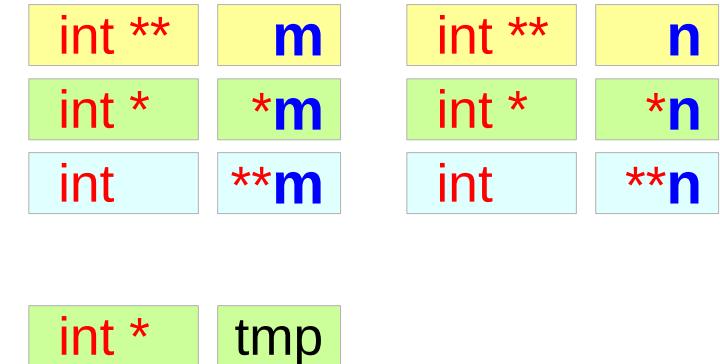
Annotations on the right side of the diagram identify the rows:

- The middle row is labeled "function call".
- The bottom row is labeled "function prototype".

# Pass by integer pointer reference

```
void pswap (int **m, int **n)
{
    int * tmp;

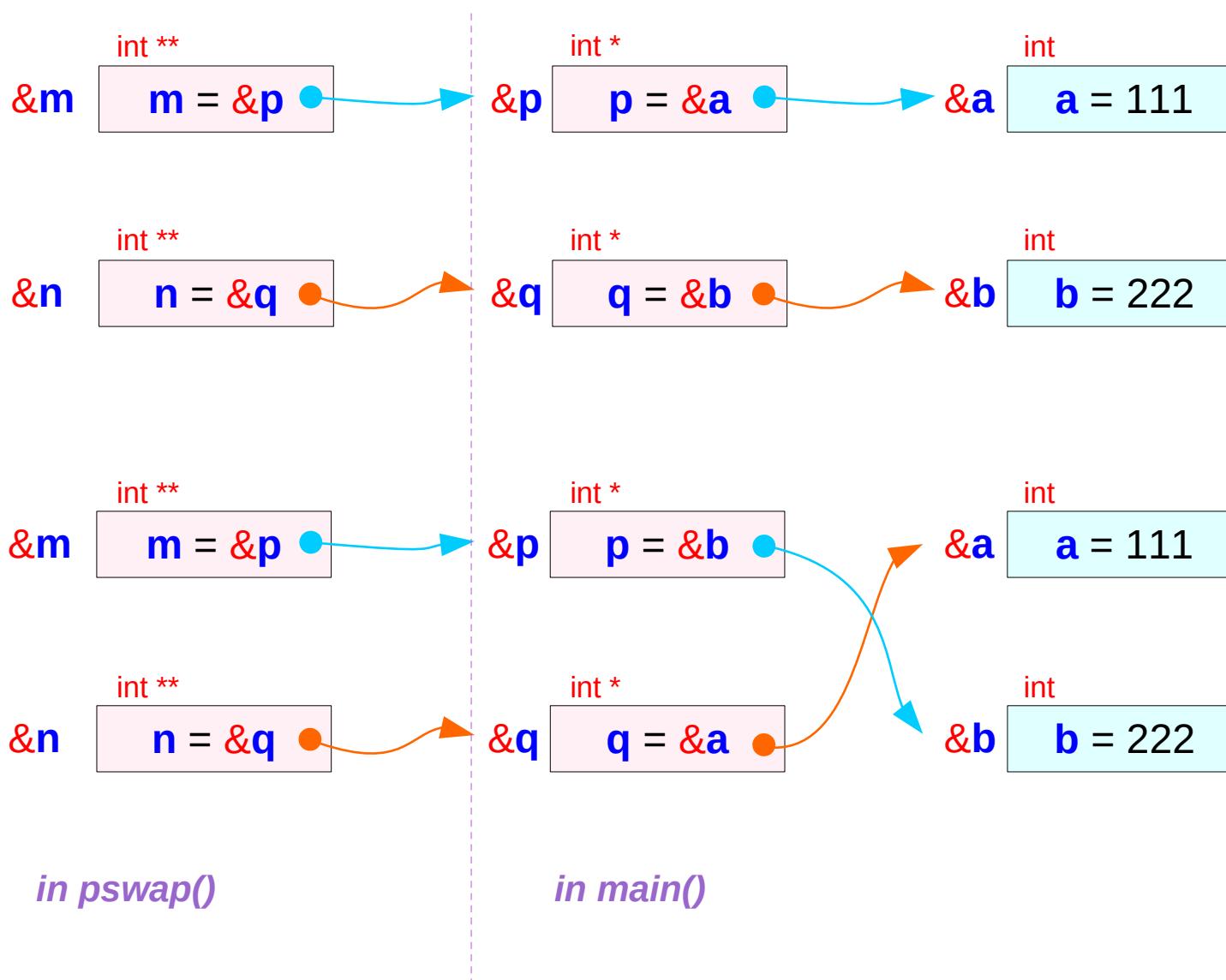
    tmp = *m;
    *m = *n;
    *n = tmp;
}
```



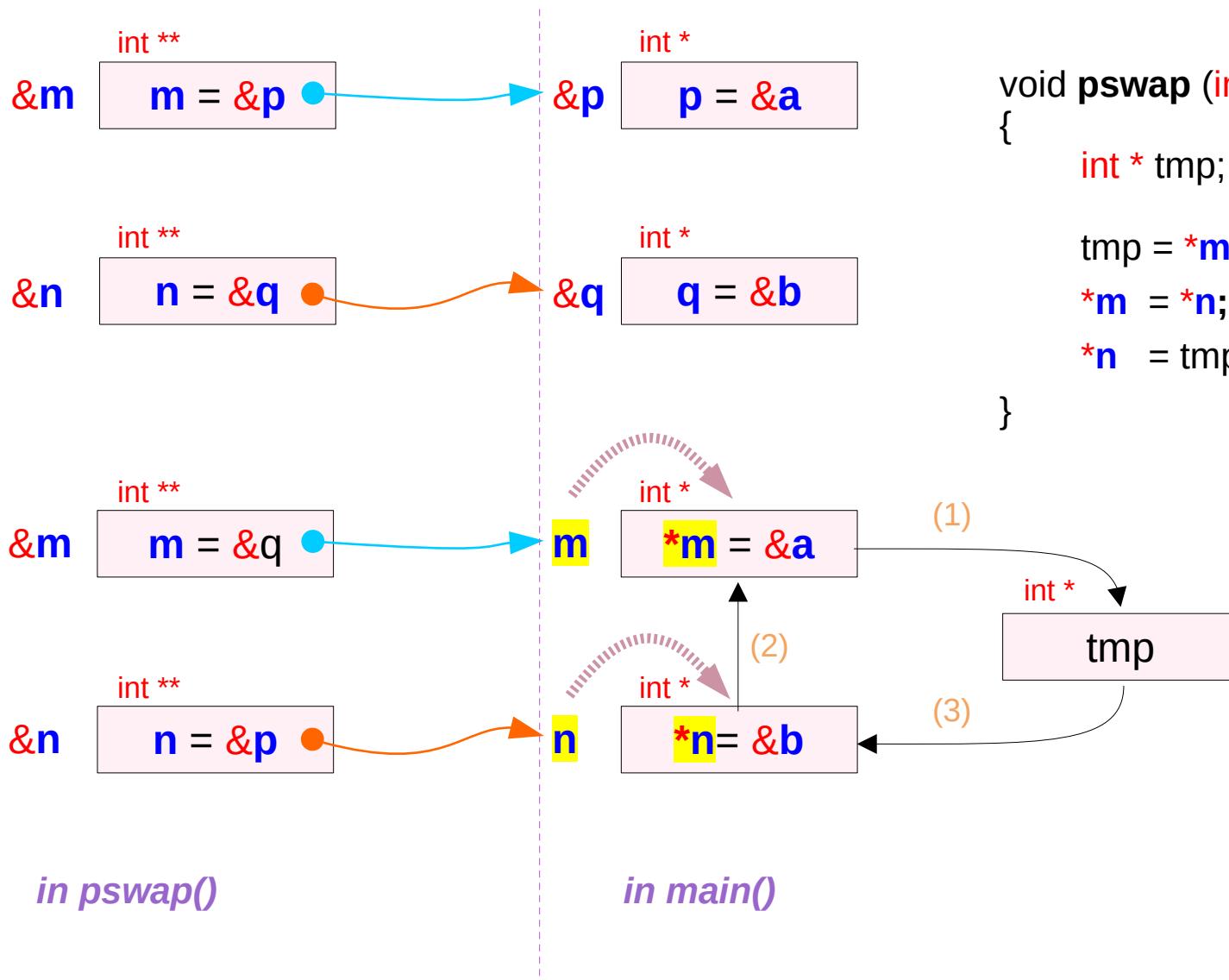
```
int a = 111, b = 222;
int *p = &a, *q = &b;
...
pswap ( &p, &q );
```

int \*\* m  
int \* \*m  
int \*\*m

# Swapping integer pointers



# Swapping integer pointers via double pointers



```
void pswap (int **m, int **n)
{
    int * tmp;

    tmp = *m;          (1)
    *m = *n;          (2)
    *n = tmp;          (3)
}
```

## References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun